# CYBER ASSURED SYSTEMS ENGINEERING WITH AADL

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UNITED STATES ARMY

### TEAM

#### COLLINS-LED DARPA CYBER ASSURED SYSTEMS ENGINEERING (CASE) PROJECT

- Collins Aerospace
  - Architectural transformations for cyber-resilience
  - Component synthesis and proofs
  - Formal analysis and assurance case
  - Tool integration
- Data 61
  - seL4 formally verified secure microkernel for memory protection
  - Formally verified components (seL4, CakeML language)
- University of Kansas
  - Formally verified attestation for distributed computing platforms
- Adventium
  - Real-time scheduling
  - AADL modeling and code generation
- Kansas State University
  - Automatic code generation from architecture models with proof of equivalence
  - Information flow analysis















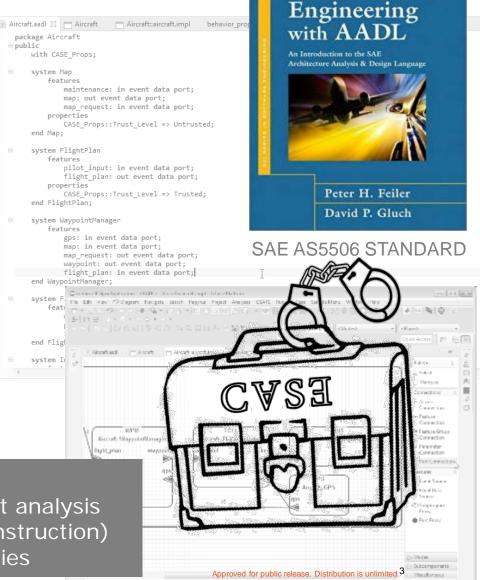
# BRIEFCASE TOOL CAPABILITIES

- Integrated model-based systems engineering tool suite based on Architecture Analysis & Design Language (AADL) models
- Transform system design to satisfy cyber-resiliency requirements
- Generate new high-assurance components from formal specifications
- Verify system design using formal methods and document evidence/compliance with assurance case
- Generate software integration code directly from verified architecture models, targeting multiple operating systems (including seL4)

#### Why AADL?

ins Aerospace

- Sufficiently rigorous semantics to support analysis
- Correct level of abstraction (supports construction)
- OSATE supports addition of new capabilities



Model-Based

### 3 TECHNOLOGY PILLARS

#### 1. Developer assistance to implement cyber-resiliency

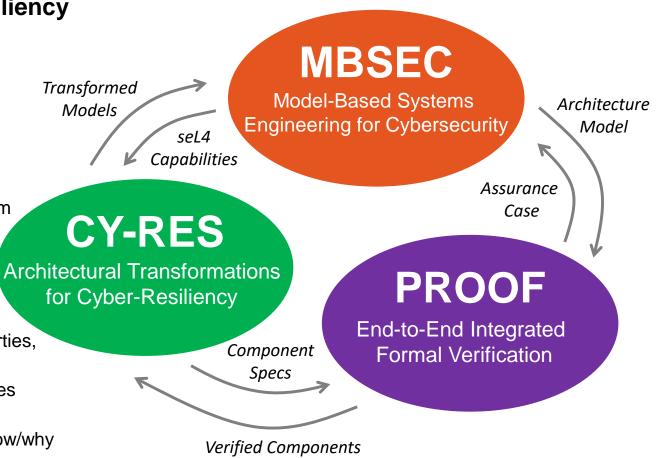
- Automated architecture transforms for threat mitigation
- High assurance components generated from specifications
- Techniques to deal with legacy code (cyber retrofit)

#### 2. MBSE environment for high-assurance cyberresilient system development

- Build system directly from detailed, verified AADL model
- Makes the security guarantees of seL4 accessible to system developers
- Ability to target different platforms to facilitate incremental debugging/development

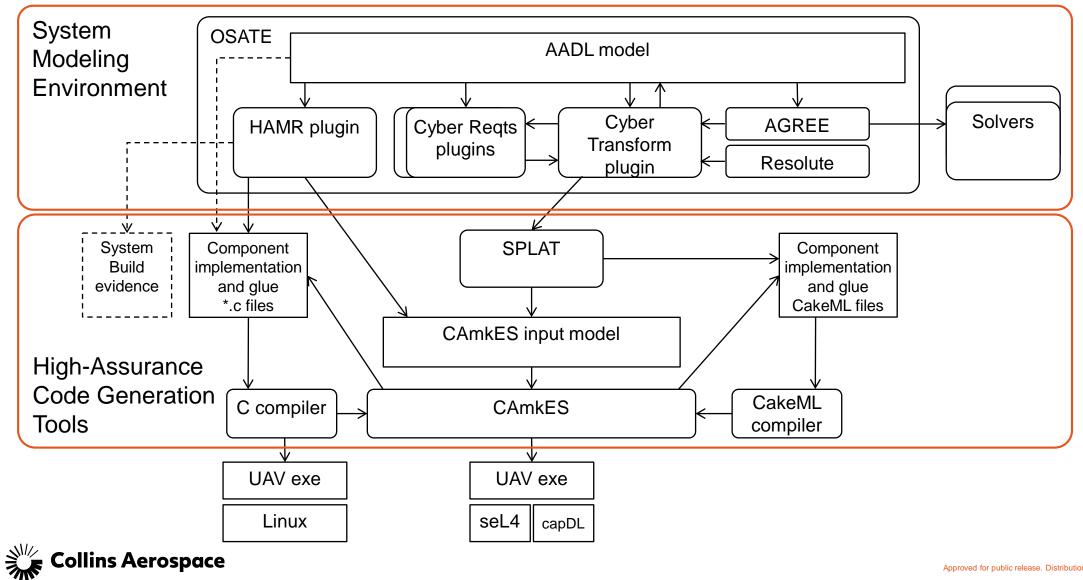
#### 3. Integration of formal verification/proof

- Formal verification of functional and cyber-resiliency properties, information flow properties, component proofs
- Code generation equivalence to model, seL4 build preserves properties
- Integrate evidence as an assurance case demonstrating how/why requirements are satisfied



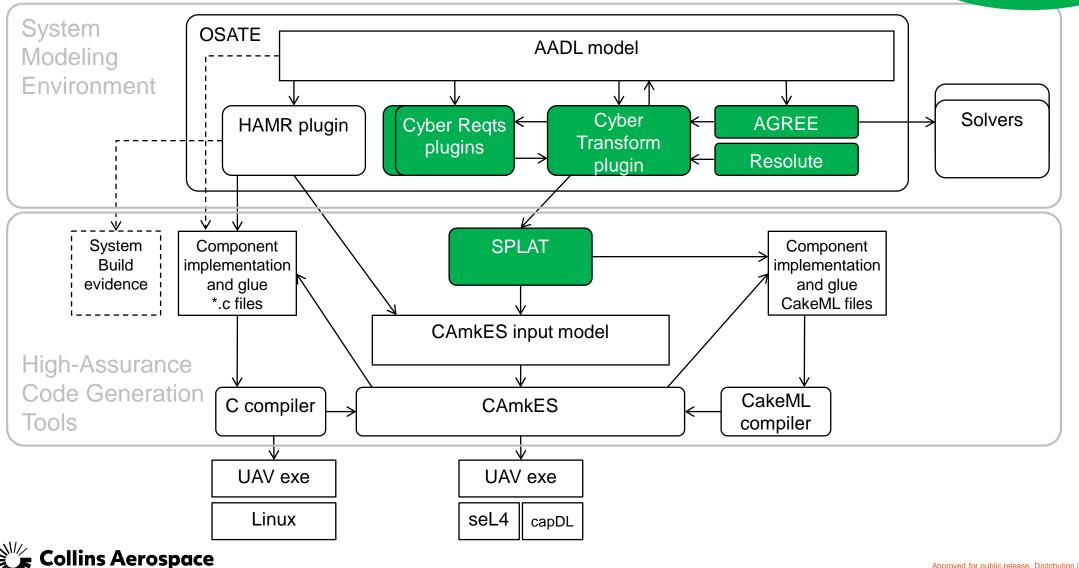


#### BRIEFCASE TOOL ARCHITECTURE



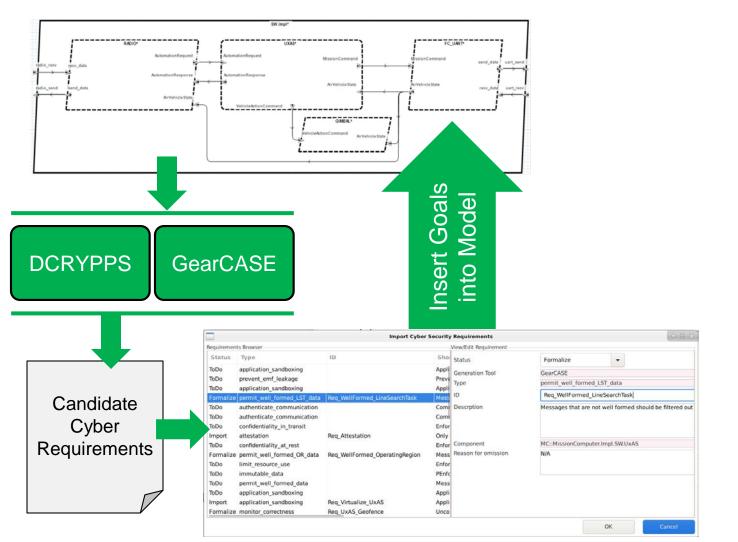
#### ARCHITECTURAL TRANSFORMATIONS FOR **CYBER-RESILIENCY**

1. CY-RES



# 1. GENERATE / IMPORT CYBER REQUIREMENTS

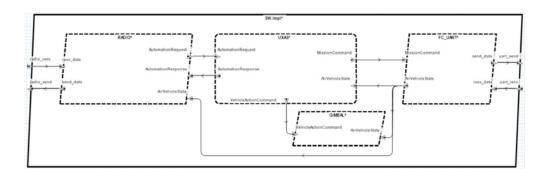
- Choose one of the Cyber Requirements generation tools
  - CRA GearCASE plugin
  - Vanderbilt/DOLL DCRYPPS plugin
- Initial model data is exported to selected tool
- Requirements import wizard manages the generated requirements
  - Select action
  - Naming/tagging
  - Associate with formal properties
- Requirements inserted into model as Resolute goals (GSN)
  - We will build an assurance argument to satisfy these goals

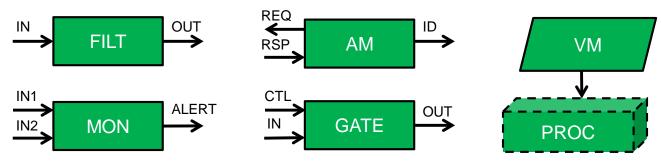


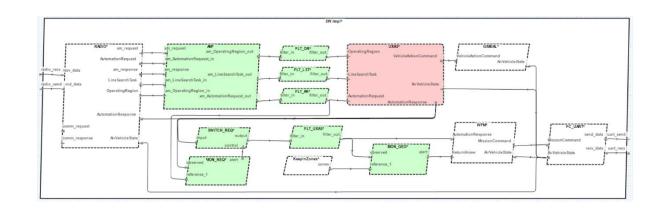


# 2. APPLY CYBER TRANSFORMATIONS

- Cyber requirements tools provide model context and sometimes suggested mitigation
- System engineer selects from available cyber-resiliency transformations
  - Filter
  - Monitor
  - Gate (controlled by monitor)
  - Attestation
  - Virtualization
  - seL4 build prep
- Wizard interface collects needed configuration data
- Tool automatically transforms AADL model
- Also adds Resolute assurance case strategy to show how the associated goal (requirement) is satisfied









# 2A. INSERT ASSURANCE CASE STRATEGY

- Resolute links cyber transform to goal as a GSN strategy
- Checks for violations/changes that impact correctness
- Collects evidence and generates assurance case

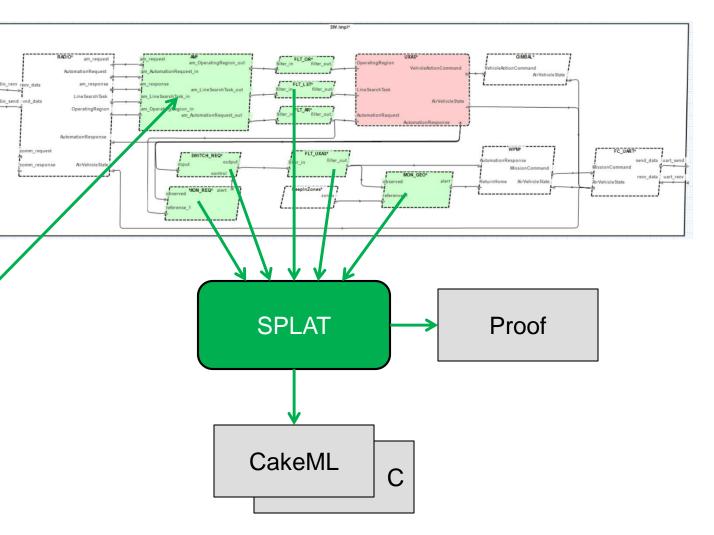




MC\_SW/impl\* Implements MC\_SW

## 3. GENERATE HIGH ASSURANCE COMPONENTS

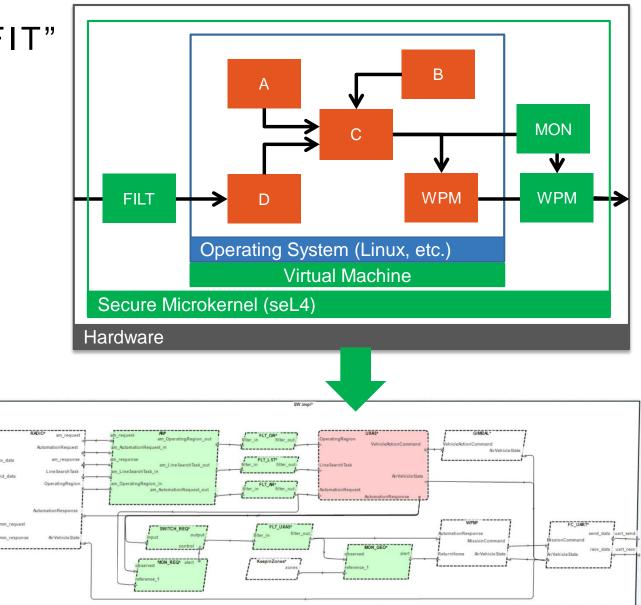
- Some of the cyber transforms insert new high-assurance components into the model
- The behavior of the component (its contract) is specified in AGREE
- SPLAT generates component implementations from their specifications
- SPLAT also generates a proof showing that the component implements its specification
- Other components (such as the Attestation Manager) are pre-built pre-verified libraries
- Their implementations are essentially library functions that are added to the build, possibly with some configuration data from the model





# USE CASE : "CYBER RETROFIT"

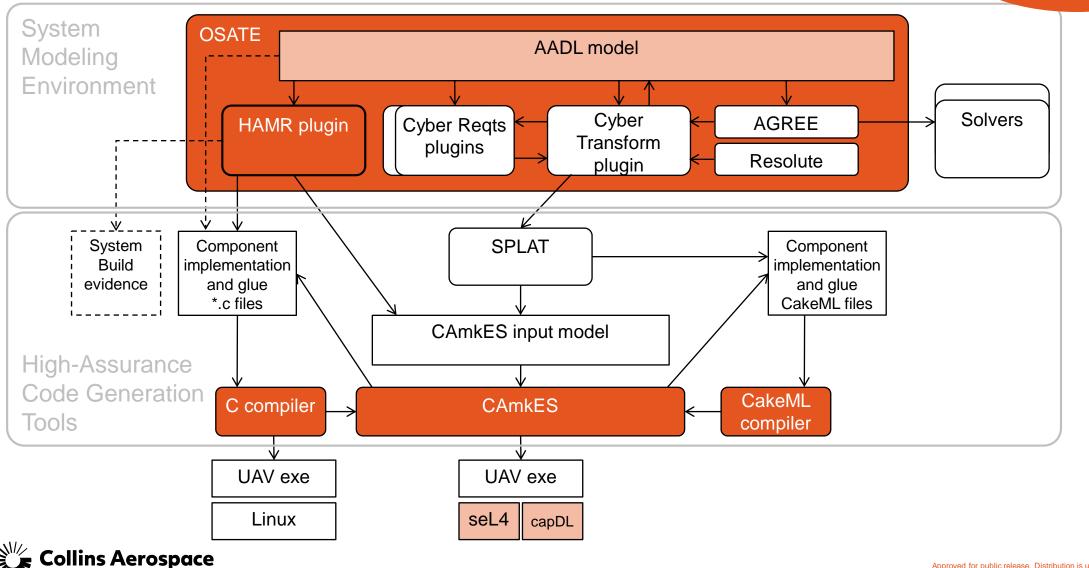
- How can we deal with legacy code/systems?
- Example: UxAS
  - Open source, less trusted
  - Not designed for cybersecurity
  - Written in C++
- 1. Virtualize legacy system
  - Support for multiple VMs, if needed
- 2. Host on seL4
- 3. Extract/harden critical components
- 4. Filter inputs
- 5. Monitor outputs





#### MODEL-BASED SYSTEMS ENGINEERING FOR CYBERSECURITY

2. MBSEC



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# MBSEC

#### MODEL-BASED SYSTEMS ENGINEERING FOR CYBERSECURITY

#### • The model IS the system

- It is not just a picture for communication, or a way to capture requirements, or a means to perform analyses
- It is all of those things, but it is also how we will actually build the system, and do so in a way that ensures compliance with the design captured in the model
- Rationale for AADL
  - Provides the necessary precision and semantics for building real-time distributed embedded software systems
- MBSE for seL4
  - Furthermore, we are bridging the gap between systems engineers and the formally verified **seL4 microkernel**
  - Ensure usability by systems engineers and promote successful transition
- HAMR
  - High Assurance Modeling for Rapid Engineering
  - Code generation from AADL models



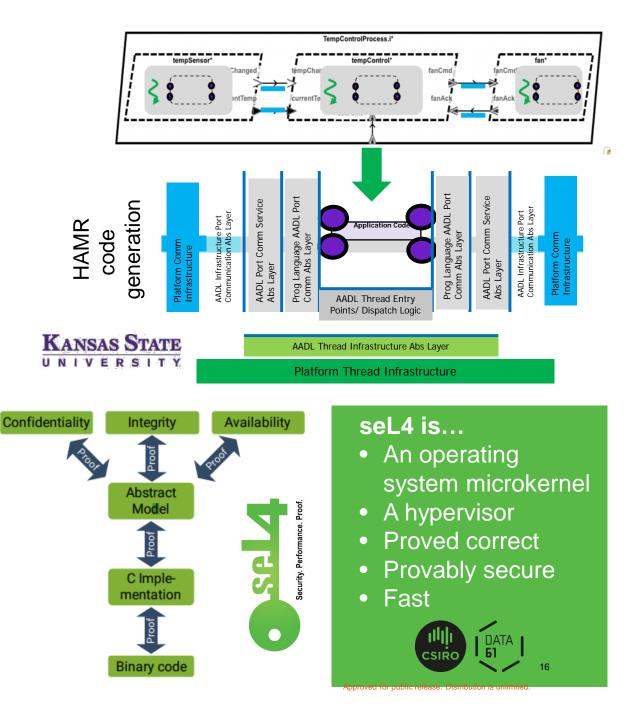


# SOFTWARE INFRASTRUCTURE

#### HAMR AND SEL4

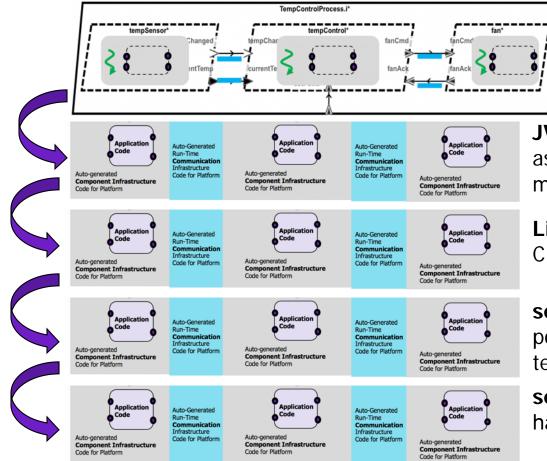
- HAMR is a multi-stage translation architecture to address CASE goals of component migration between platforms and information flow control
- Semantic consistency from model to execution
- Ensures model-level analysis applies to deployed code
- Same computational model across different platforms
- Build for multiple target platforms:
  - seL4 / Linux / Virtual Machine
  - Build for workstation / emulator / embedded platform
- seL4 microkernel guarantees partitioning of components and communication, backed by computer-checked proofs
- seL4 guarantees no infiltration, exfiltration, eavesdropping, interference, and provides fault containment for untrusted code





#### HAMR SUPPORTS MULTIPLE LANGUAGE/ PLATFORM COMBINATIONS

The flexibility of being able to easily shift between different platforms was quite useful as the team experimented with building the Phase 2 Experimental Platform assessment deliverable.



**AADL / OSATE** – design model, types, perform analyses

**JVM/Slang** – data types, port constraints, basic aspects of application logic, initial unit testing – some mocked up components, many useful visualizations

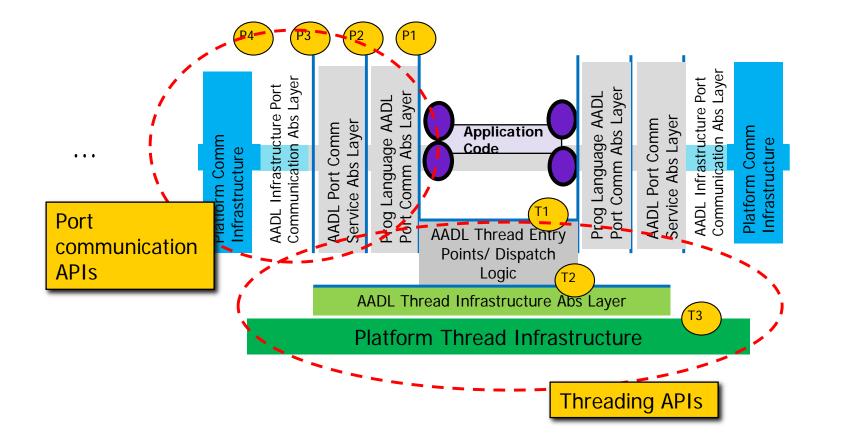
**Linux C** – compile Slang to C, or manually code C, and debug C implementation, VMs mocked up

**seL4 C / Qemu** – C application code easily ports to seL4 native components, add in VMs, test/simulate/debug in Qemu

**seL4 C / board** – seL4 build shifted to actual hardware for final testing and deployment

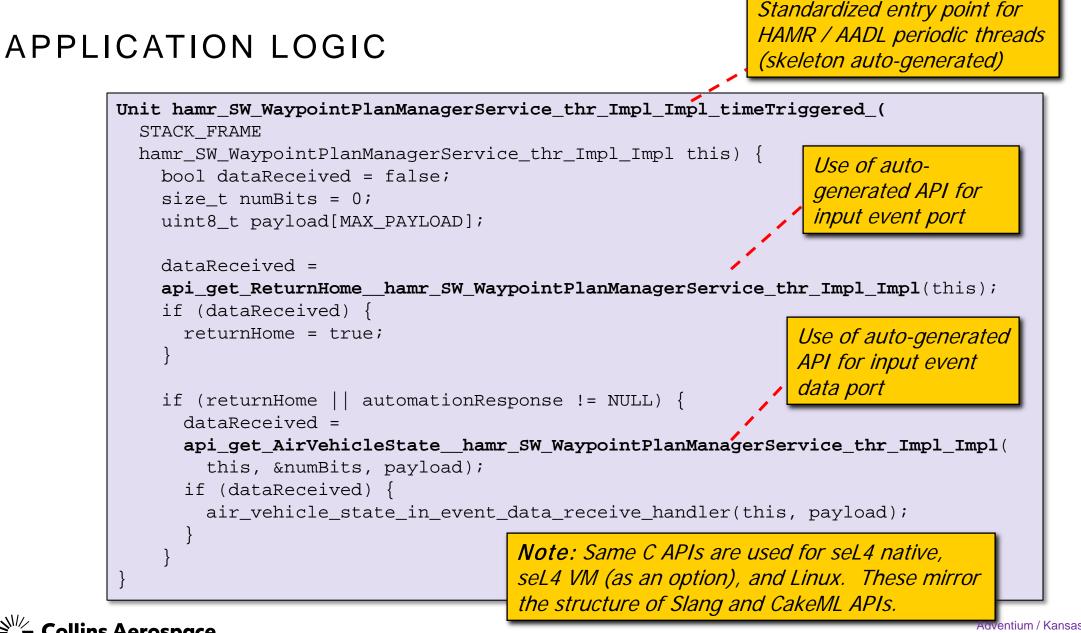


#### HAMR ABSTRACTION LAYERS





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#### **Collins Aerospace**

# RESOLINT: LINTER TOOL FOR AADL MODELS

dispatch\_pro

valid dispate

Description

Service Stress (2 items)

> 🙆 Warnings (54 items)

- Define *rules* in Resolute that correspond to modeling guidelines
- Group rules into *rulesets* corresponding to organizational process, customer requirements, certification guidelines, and tool constraints
- Automatically check compliance with modeling guidelines in OSATE

ol_specified	Threads should have the dispatch_protocol property specified	٨	<pre>** "Threads should have the Dispatch_Protocol property specified" ** forall (t : thread) .     lint_check(t, has_property(t, Dispatch_Protocol)) valid_dispatch_protocol() &lt;=</pre>				
protocol	Threads can only specify a dispatch_protocol property of <i>periodic</i> or <i>sporadic</i>	0	**	<pre>** "Threads can only specify a dispatch_protocol property of periodic or sporadic" ** foral1(t : thread) . lint_check(t, has_property(t, Dispatch_Protocol) =&gt;     (property(t, Dispatch_Protocol) = "Sporadic" or property(t, Dispatch_Protocol) = "Periodic" or protocol = "Periodic" or protocol = "Periodic" or protocol = "Periodic" or protocol = "Peri</pre>			
				2 3 <sup>©</sup> public	<pre>ge System_Build c nnex resolute {**    ruleset HAMR {         info (print("Linting HAMR ruleset"))         error (one process()))</pre>		
				11 12 13 14 15 16 17 18 19 20 21 22	<pre>warning (dispatch_protocol_specified()) error (valid_dispatch_protocol()) error (bounded_integers()) error (bounded_floats()) warning (data_type_specified()) warning (subcomponent_type_specified()) error (array_dimension()) error (one dimensional arrays())</pre>		

~

Only one processor-bound process can contain thread or thread-group subcomponents

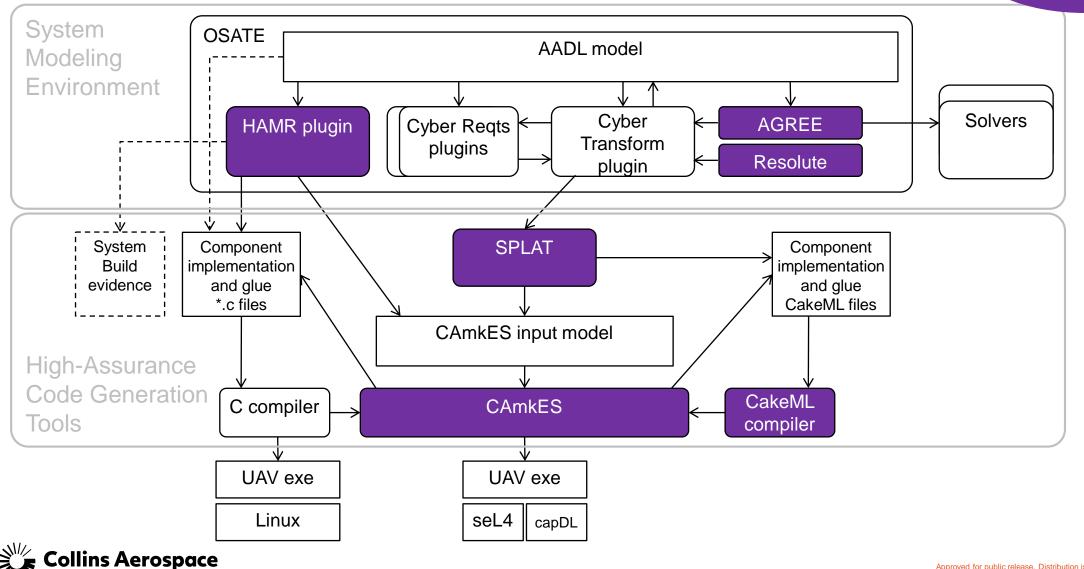
Ocmponents bound to a virtual processor may only communicate with components bound to other processors via event data ports

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#### END-TO-END INTEGRATED FORMAL VERIFICATION

3. PROOF



#### END-TO-END INTEGRATED FORMAL VERIFICATION

system pro		
architecture	assurance case	
high-assurance components	legacy components	
HAMR correspo		
CAmkES trans		
seL4 initializ		
seL4 p		



#### BRIDGING THE GAP WITH PROOFS

SP

seL4 integrity policy (access control) is enforced on a running system seL4 system initializer constructs system whose capabilities and objects conform to capDL specification

translates system HAMR translates 3 CAmkES description into capDL spec

CAmkES system description

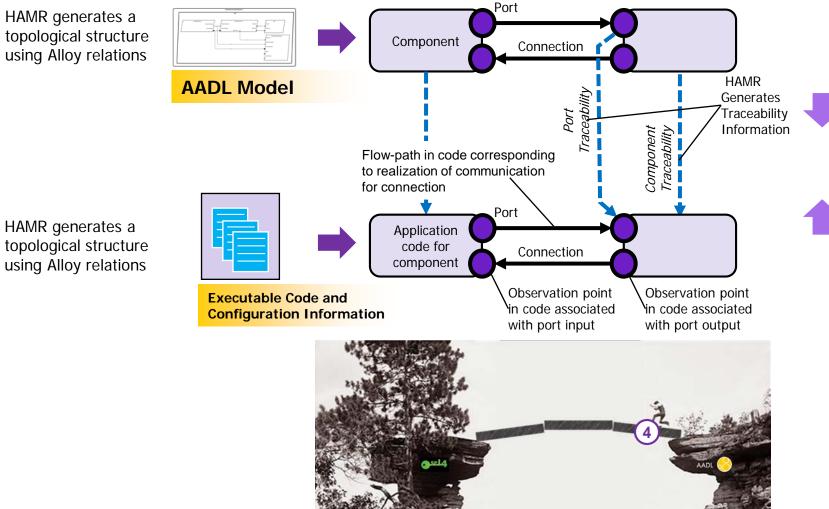
**Collins Aerospace** 

AADL

# HAMR CORRESPONDENCE PROOF CONCEPT (4)

HAMR generates a topological structure using Alloy relations

HAMR generates a



FlowPreservation (formal spec in Alloy): For every connection between two components in AADL, there is a flow path in the source code between code artifacts associated with the ports.

**NoNewFlows** (formal spec in Alloy): For every flow path between two components in the source code, there is a connection in the AADL model between corresponding ports.

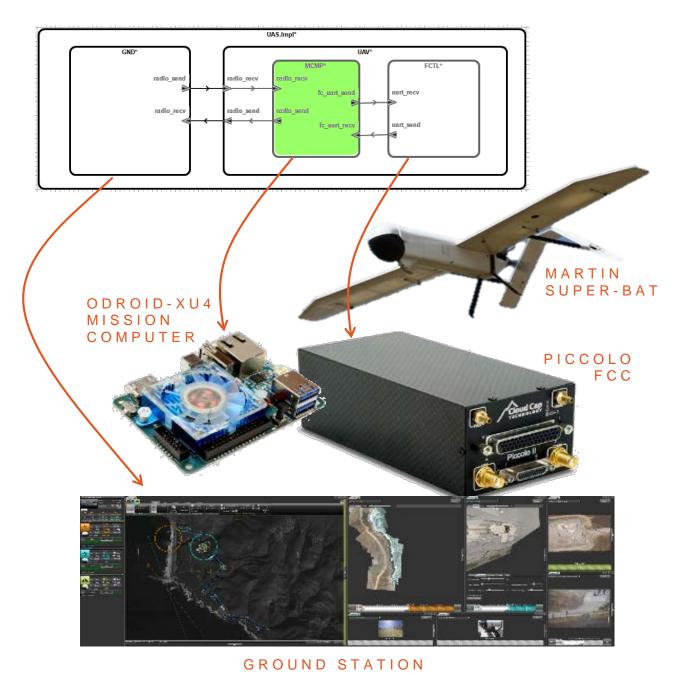


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#### PHASE 2 UAV DEMONSTRATION

#### AFRL UXAS AUTONOMOUS MISSION PLANNER

- Ground Station sends automation requests to UAV
- UAV Mission Computer processes requests and generates flight plans
- UAV Flight Control Computer computes guidance commands to follow current segment of flight plan
- CASE evaluation team developed cyber attacks on baseline platform
- Collins team hardened platform using BriefCASE tools
- Evaluation team attacks ineffective against hardened platform

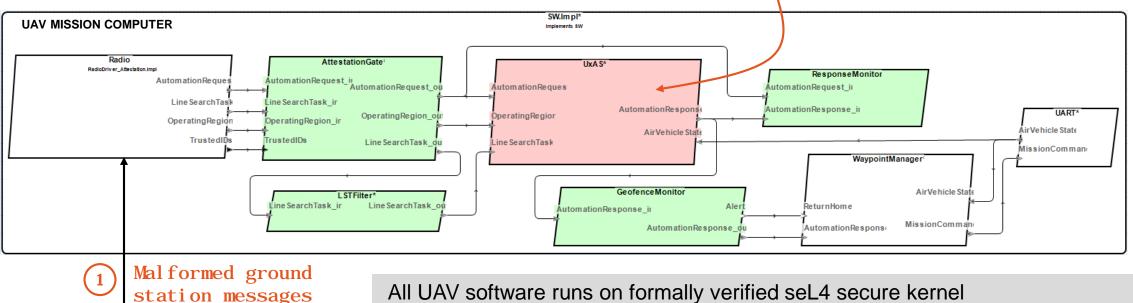




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# ATTACKS / MITIGATION

2 Troj ans added to UxAS PlanBuilderService Replaces AutomationResponse with plan that violates KeepOutZone



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#### GND STN

3 Code added to ground station software to generate malicious requests

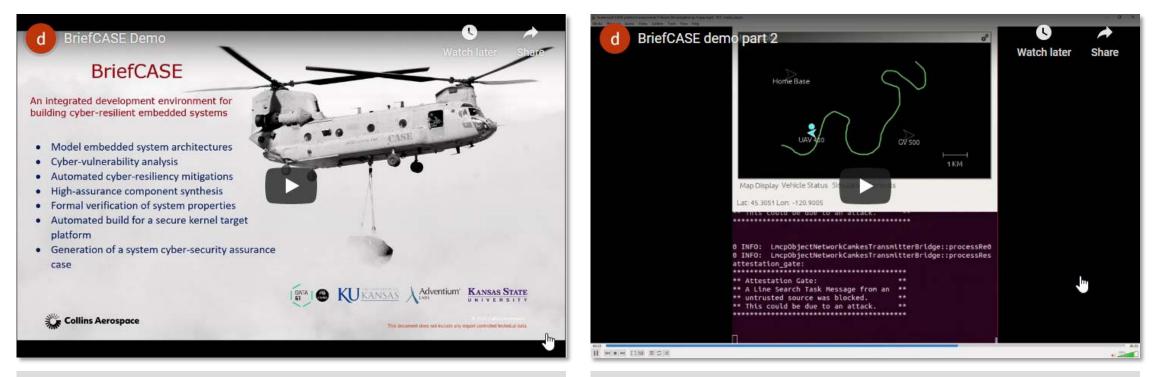


All UAV software runs on formally verified seL4 secure kernel UxAS software isolated in Linux virtual machine (cyber retrofit) Specific attacks from evaluation team mitigated:

- 1. Malformed messages blocked by high-assurance filter
- High-assurance geofence monitor detects plan that violates KeepOutZone and sends alert to WaypointManager to trigger return to base Response monitor detects crashed UxAS planner and alerts operator
- 3. Remote attestation measures ground station software and detects modified code

# CASE PHASE 2 DEMO VIDEOS

#### HTTP://LOONWERKS.COM/PROJECTS/CASE



This video provides a demonstration of the BriefCASE tool environment, showing how to use the tools to address multiple cyber-resiliency requirements for a UAV mission computing system (22:35).

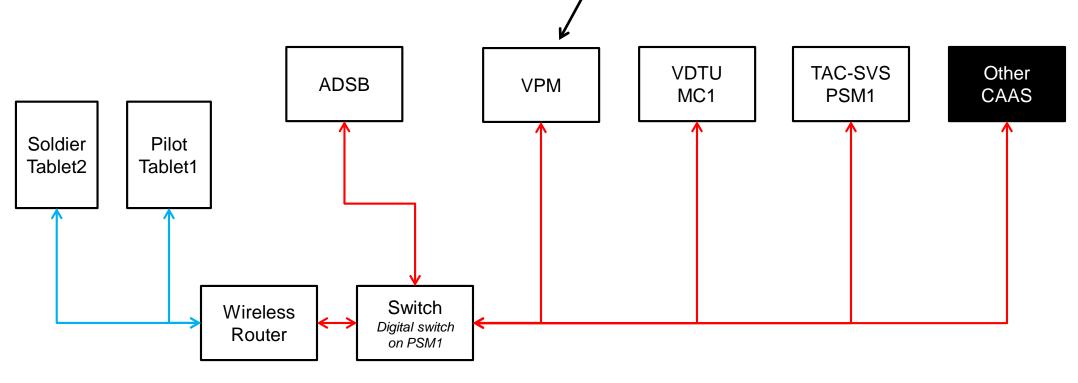
In part two, we run the hardened UAV mission computing system built in the first video and test it against several cyber attacks to show the effectiveness of the approach (10:13).



### CASE PHASE 3 DEMONSTRATION

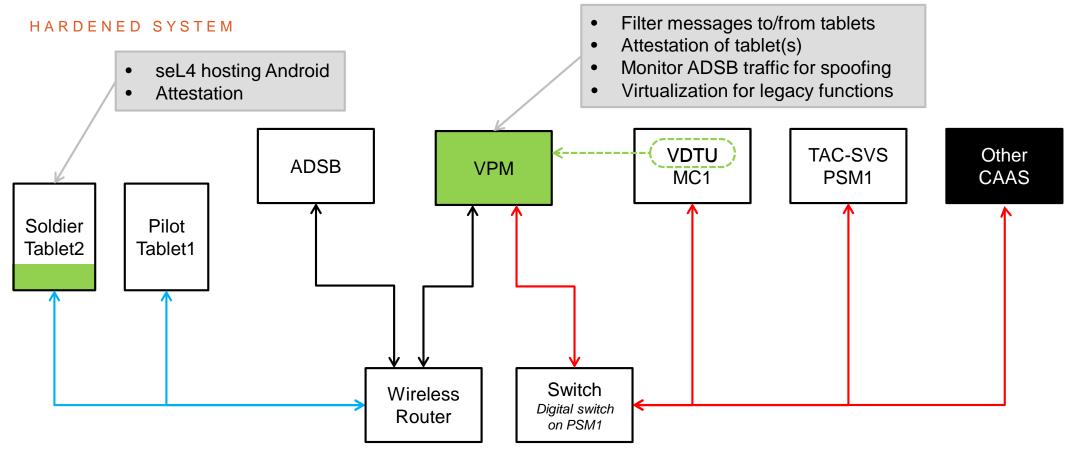
#### BASELINE SYSTEM (PHYSICAL)

- Collins Common Avionics Architecture System (CAAS)
- Goal : Extend (securely) to add wireless connectivity
- Idea : Repurpose one of the mission processing modules (VPM) to act as guard using BriefCASE AADL tools





### CASE PHASE 3 DEMONSTRATION





#### CONCLUSION

#### Developer assistance to implement cyber-resiliency

- Automated architecture transforms for threat mitigation
- High assurance components generated from specifications
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Open-source tools : GitHub.com/Loonwerks/formal-methods-workbench Demo video : Loonwerk.com/projects/case.html

#### **CY-RES**

Architectural Transformations for Cyber-Resiliency

#### **MBSEC**

Model-Based Systems Engineering for Cybersecurity

#### PROOF

End-to-End Integrated Formal Verification

