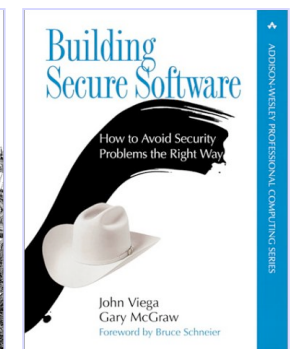
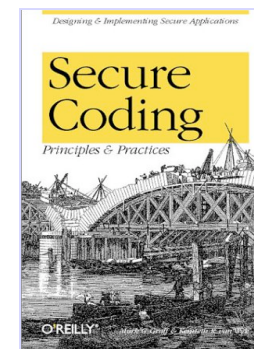
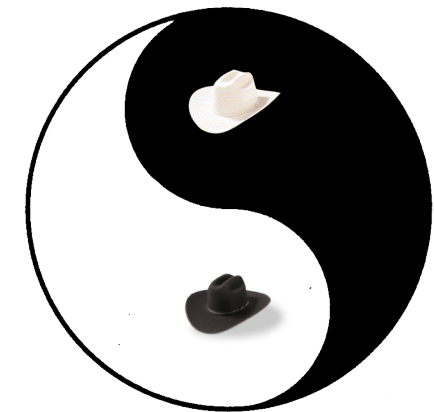




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Setting the stage





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Agenda

- | | |
|-------------|----------------------------------|
| 9:00-10:00 | Software [in]security |
| 10:15-12:00 | Exploiting Software and exercise |
| 1:00-2:30 | Software security touchpoints |
| 2:45-4:30 | Seven pernicious kingdoms |
| 4:30-5:00 | Code review and next steps |



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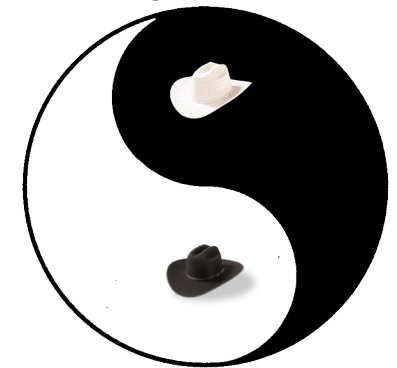
Pop quiz

- What do wireless devices, cell phones, PDAs, browsers, operating systems, servers, personal computers, routers, public key infrastructure systems, and firewalls have in common?



Questions for you

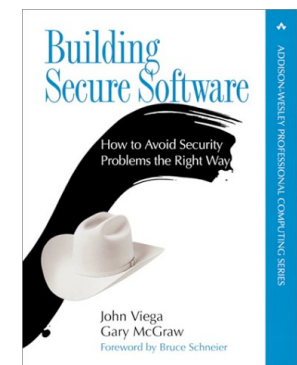
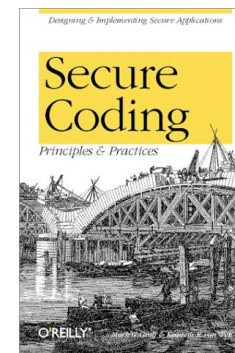
- Who is from dev? How about testing? Anyone here from product management?
- What languages do you use? C? C++? Java?
- How do you describe and capture software architecture and design?
- Do you follow a particular software process in your group?





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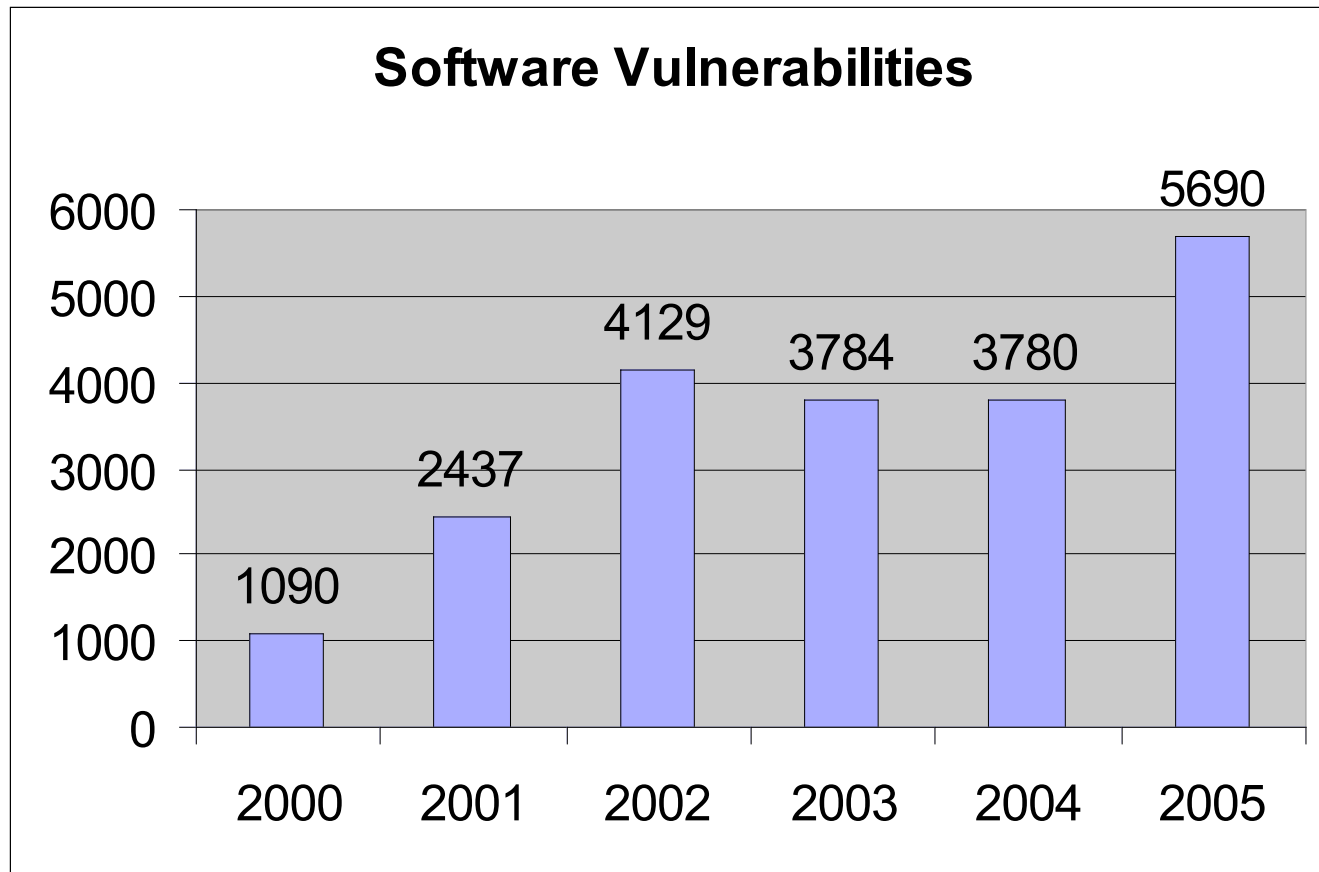
Software [in]security



The Problem



Software vulnerability growth





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The Trinity Of Trouble: Connectivity

- The Internet is everywhere and most of our software is on it
- When was the last time that you did business with a major vendor who had no Internet connectivity?
- Tried VoIP on your mobile phone in a coffee shop WiFi hotspot yet?

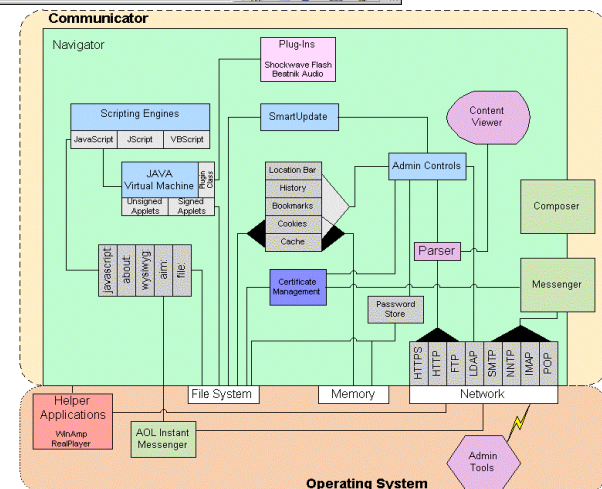
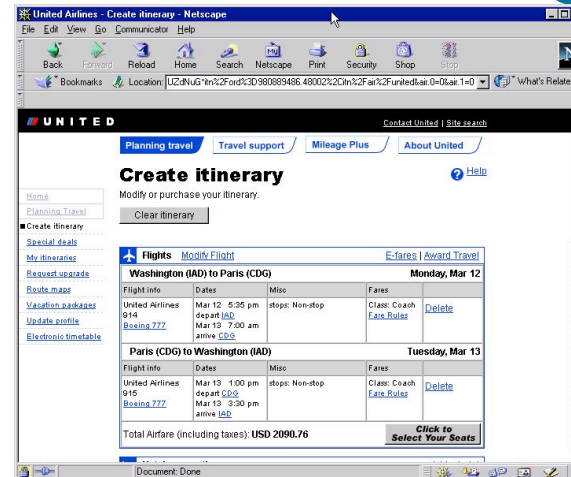
The network is
the computer.





The Trinity Of Trouble: Complexity

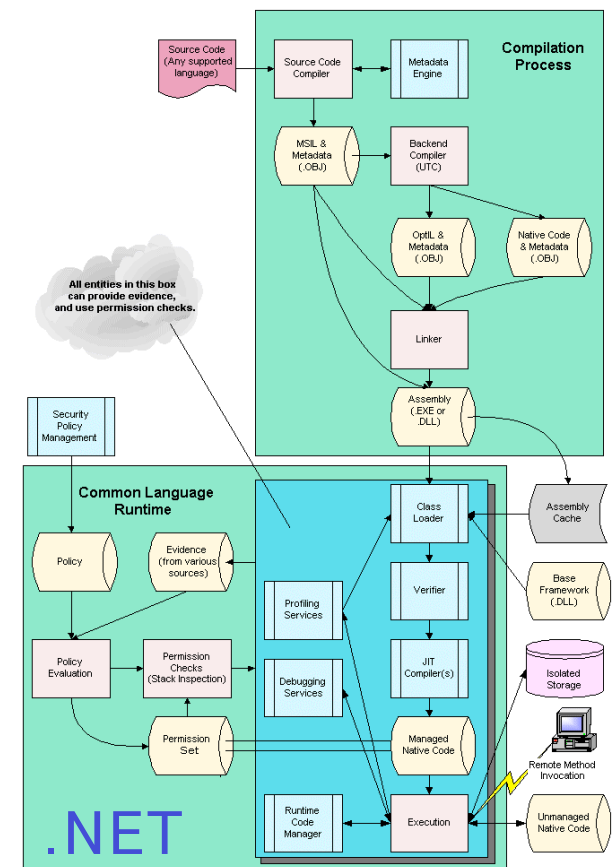
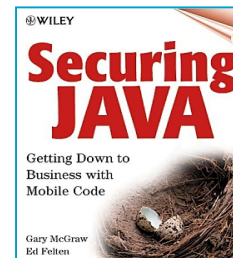
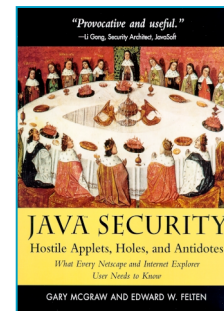
- A simple user interface can be enormously complex “under the hood”
- Consider what happens behind the scenes in one of today’s AJAX web applications
- But it sure does make for a compelling “user experience”





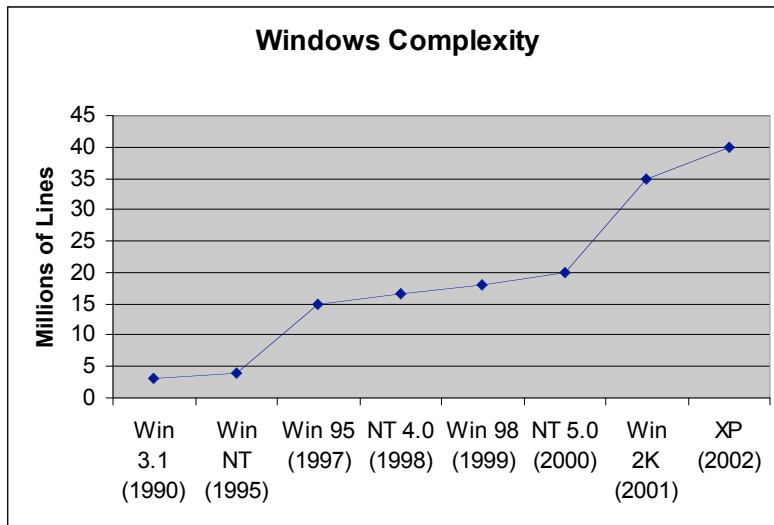
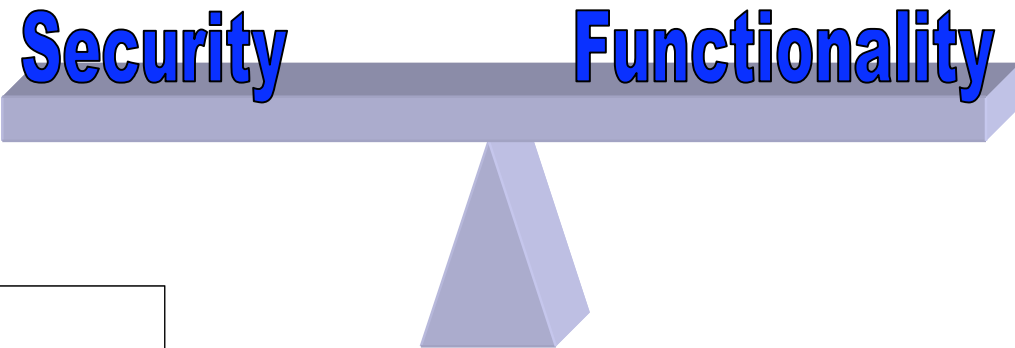
The Trinity Of Trouble: Extensibility

- Systems evolve in unexpected ways and are changed on the fly
- After all, who would want a computing device that can't be functionally extended?
- From J2ME to desktop PC users (running with administrative privileges)





The classic security tradeoff





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So what's the problem?

- Well, for starters
 - Consumers don't demand more
 - Software developers tend to lack knowledge of vulnerabilities, attacks, and threats
 - IT security tends to not understand software development
- But that's not all!



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Additional problems - 1

- We don't pay enough attention to our failures
- Consider other engineering disciplines





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Additional problems - 2

- We fail to consider business risks first and foremost
- Business must drive technology
- Consider Wi-Fi, Word macros, USB drives, etc.





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Additional problems - 3

- Old school information security solutions don't adequately protect the software
- Consider IM, Skype, Wi-Fi, VPNs





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Additional problems - 4

- Software testing does not adequately address security
- Penetration testing is not sufficient





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Additional problems - 5

- Too much attention is paid to functional spec
- Consider what can go wrong as well





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Additional problems - 6

- IT security is viewed as an impediment to business
- Don't just be the person that says no

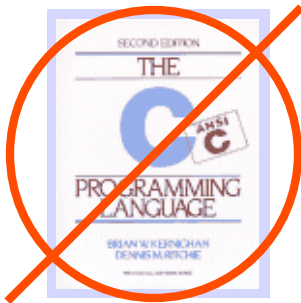




Security problems are complicated

IMPLEMENTATION BUGS

- Buffer overflow
 - String format
 - One-stage attacks
- Race conditions
 - TOCTOU (time of check to time of use)
- Unsafe environment variables
- Unsafe system calls
 - System()
- Untrusted input problems



ARCHITECTURAL FLAWS

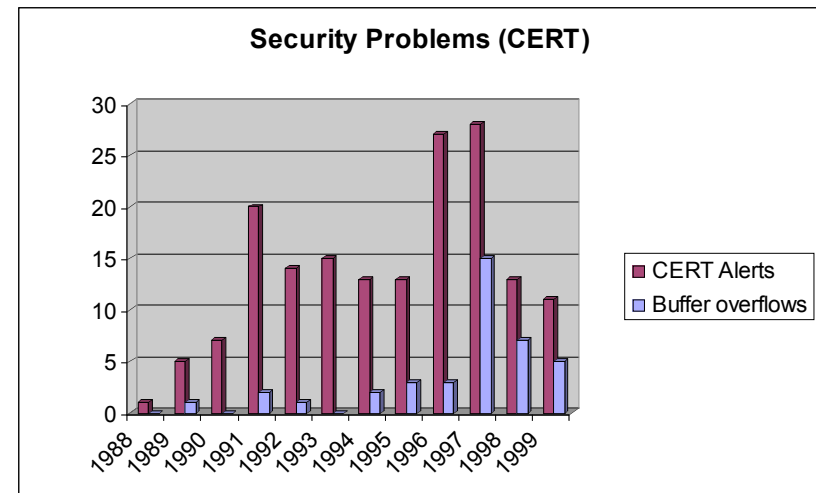
- Misuse of cryptography
- Compartmentalization problems in design
- Privileged block protection failure (DoPrivilege())
- Catastrophic security failure (fragility)
- Type safety confusion error
- Insecure auditing
- Broken or illogical access control (RBAC over tiers)
- Method over-riding problems (subclass issues)
- Signing too much code



BUG: The dreaded buffer overflow

- Overwriting the bounds of data objects
- Allocate some bytes, but the language doesn't care if you try to use more
 - `char x[12];`
`x[12] = '\\0';`
- Why was this done? Efficiency!
- Two main flavors of buffers
 - Heap allocated buffers
 - Stack allocated buffers
 - Smashing the stack is the most common attack

- The most pervasive security problem today in terms of reported bugs





Pervasive C problems lead to BUGS

```
void main() {  
    char buf[1024];  
    gets(buf);  
}
```

- How not to get input
 - Attacker can send an infinite string!



Chapter 7 of
K&R (page 164)

- Calls to watch out for

Instead of:	Use:
gets(buf)	fgets(buf, size, stdin)
strcpy(dst, src)	strncpy(dst, src, n)
strcat(dst, src)	strncat(dst, src, n)
sprintf(buf, fmt, a1,...)	snprintf(buf, fmt, a1, n1,...) (where available)
*scanf(...)	Your own parsing

- Hundreds of such calls
- Use static analysis to find these problems
 - ITS4, Fortify
- Careful code review is necessary



FLAW: 802.11b WEP crypto

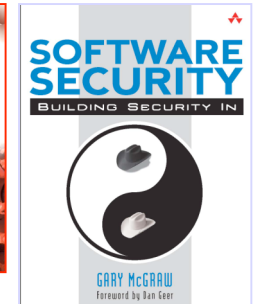
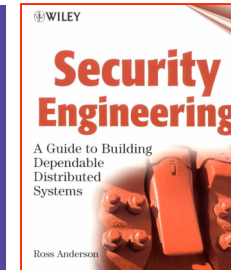
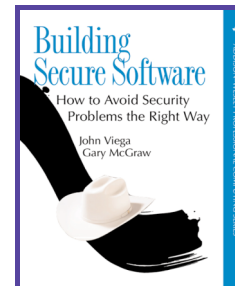
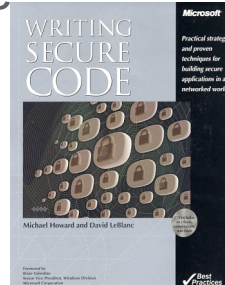
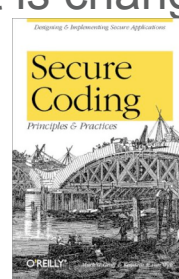
- Well-documented flaws in the design of the WEP protocol
- Even if implemented 100% perfectly, the design is flawed and the encryption easily circumvented
- 802.11b is widely deployed and wildly popular
- It was designed by experts
- Would you entrust a mission-critical enterprise to run over it?





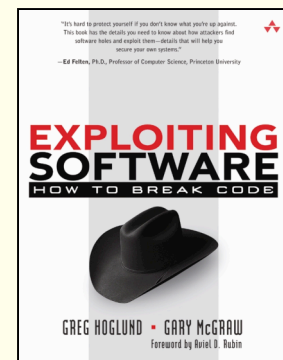
Software security: state of the practice

- Programming is hard
- Popular languages are really awful (C/C++)
- Many subtleties to learn
- Lots to know
- The only constant is change
- Some good resources on software security
- Tools are getting better, but only cover BUGS



Software security is not security software!
Software security is about building things properly.

Exploiting software

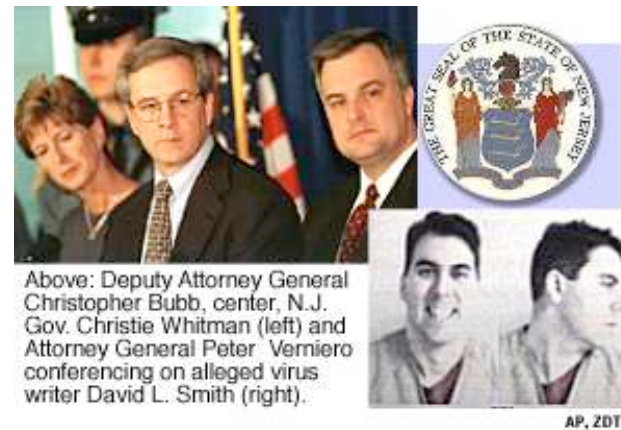




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Who is the bad guy?

- Hackers
 - “Full disclosure” zealots
- “Script kiddies”
- Criminals
 - Lone guns or organized
- Malicious insiders
 - Compiler wielders
- Business competition
- Police, press, terrorists, intelligence agencies





Attackers do not distinguish bugs and flaws

- Both bugs and flaws lead to vulnerabilities that can be exploited
- Attackers are pragmatic in their approach
- Attackers write code to break your software's design and/or implementation





How attacks unfold

- Attacking a system is a process of discovery and exploration
 - Qualify target (focus on input points)
 - Determine what transactions the input points allow
 - **Apply relevant attack patterns**
 - **Cycle through observation loop**
 - **Find vulnerability**
 - **Build an exploit**

The standard process

- Scan network
- Build a network map
- Pick target system
- Identify OS stack
- Port scan
- Determine target components
- Choose attack patterns
- Leverage environment faults
- Use indirection
- Plant backdoor



Attacker's toolkit: disassemblers and decompilers

- Source code is not a necessity for software exploit
- Binary is just as easy to understand as source code
- Disassemblers and decompilers are essential tools
- Reverse engineering is common and must be understood (not outlawed)
- IDA allows plugins to be created
- Use bulk auditing

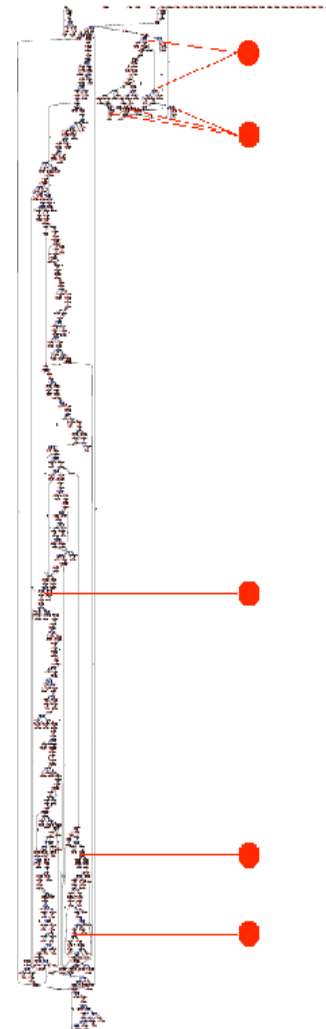


IDA Pro
by Ilfak Guilfanov



Attacker's toolkit: control flow and coverage

- Tracing input as it flows through software is an excellent method
- Exploiting differences between versions is also common
- Code coverage tools help you know where you have gotten in a program
 - dyninstAPI (Maryland)
 - Figure out how to get to particular system calls
 - Look for data in shared buffers





Attacker's toolkit: APISPY32

- Look for broken system calls (at all levels in code)
- Istrcpy() makes a great example
- On win32 systems, use APISPY to determine which APIs are being used by a target program
- Interposition attacks are a great thing to think about at this level

APISpy32 - http://www.internals.com

Process	PID	API
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcc8d8:"", PSTR:0xd2e730:"OpenTapiPerformanceData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcc8d8
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcc8f0:"", PSTR:0xd2e834:"CollectTapiPerformanceData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcc8f0
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcc910:"", PSTR:0xd2e62c:"CloseTapiPerformanceData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcc910
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd138:"", PSTR:0xd2e730:"OpenTcpIpPerformanceData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd138
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd158:"", PSTR:0xd2e834:"CollectTcpIpPerformanceData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd158
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd178:"", PSTR:0xd2e62c:"CloseTcpIpPerformanceData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd178
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd690:"", PSTR:0xd2e730:"OpenT5Object")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd690
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd6a0:"", PSTR:0xd2e834:"CollectT5ObjectData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd6a0
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd6b8:"", PSTR:0xd2e62c:"CloseT5Object")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd6b8
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd8e8:"", PSTR:0xd2e730:"WmiOpenPerfData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd8e8
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd8f8:"", PSTR:0xd2e834:"WmiCollectPerfData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd8f8
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xcd910:"", PSTR:0xd2e62c:"WmiClosePerfData")
SQLSERVER	0x00000b18	Istrcpy returned 0xcd910
SQLSERVER	0x00000b18	Istrcpy(PSTR:0x19feb9c:"P", PSTR:0x71ab7bec:"WinSock 2.0")
SQLSERVER	0x00000b18	Istrcpy returned 0x19feb9c
SQLSERVER	0x00000b18	Istrcpy(PSTR:0x19fedec:"P", PSTR:0x71ab7be4:"Running")
SQLSERVER	0x00000b18	Istrcpy returned 0x19fedec
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xdd7bc:"", PSTR:0xdd8c8:"%SystemRoot%\system32\mswsock.dll")
SQLSERVER	0x00000b18	Istrcpy returned 0xdd7bc
SQLSERVER	0x00000b18	Istrcpy(PSTR:0xddecc:"", PSTR:0xdd8c8:"%SystemRoot%\system32\mswsock.dll")
SQLSERVER	0x00000b18	Istrcpy returned 0xddecc
SQLSERVER	0x00000b18	Istrcov(PSTR:0xde25c:"", PSTR:0xdd8c8:"%SystemRoot%\system32\mswsock.dll")



Attacker's toolkit: breakpoints

- Breakpoints are central to debuggers
 - Use interrupt 3 on x86 architectures
- Mark entire blocks for access
- Single step at breakpoint (also as in debugging)
- Check out “The PIT” <http://www.hbgary.com>

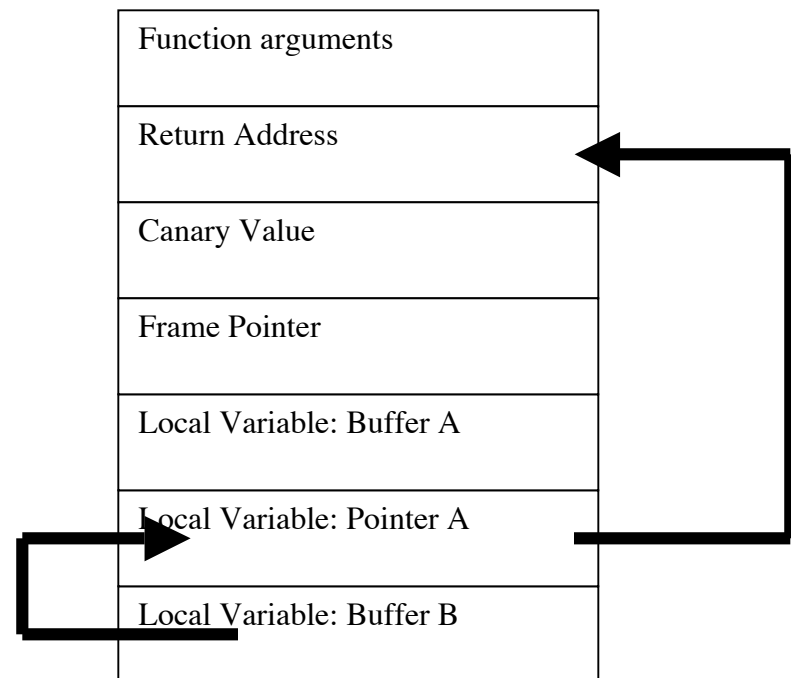




Attacker's toolkit: the buffer overflow

- Find targets with static analysis
- Change program control flow
 - Heap attacks
 - Stack smashing
 - Trampolining
- Particular examples
 - Overflow binary resource files (used against Netscape)
 - Overflow variables and tags (Yamaha MidiPlug)
 - MIME conversion fun (Sendmail)
 - HTTP cookies (apache)

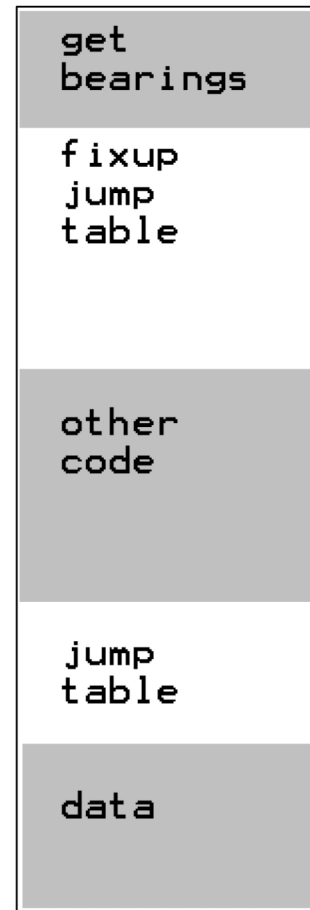
- Trampolining past a canary





Attacker's toolkit: shell code and other payloads

- Common payloads in buffer overflow attacks
- Size matters (small is critical)
- Avoid zeros
- XOR protection (also simple crypto)
- Payloads for
 - X86 (win32)
 - RISC (MIPS and sparc)
 - Multiplatform payloads





Attacker's toolkit: rootkits

- The apex of software exploit...complete the machine
- Live in the kernel
 - XP kernel rootkit in the book
 - See <http://www.rootkit.com>
- Get into the microchips (hardware viruses)
- Hide files and directories by controlling access to process tables
- Provide control and access over the network





Attacker's toolkit: other miscellaneous tools

- Debuggers (user-mode)
- Kernel debuggers
 - SoftIce
- Fault injection tools
 - Failure simulation tool
 - Hailstorm
 - Holodeck
- Boron tagging
- The “depends” tool
- Grammar rewriters



Attack Patterns





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Knowledge: 48 Attack Patterns

- Make the Client Invisible
- Target Programs That Write to Privileged OS Resources
- Use a User-Supplied Configuration File to Run Commands That Elevate Privilege
- Make Use of Configuration File Search Paths
- Direct Access to Executable Files
- Embedding Scripts within Scripts
- Leverage Executable Code in Nonexecutable Files
- Argument Injection
- Command Delimiters
- Multiple Parsers and Double Escapes
- User-Supplied Variable Passed to File System Calls
- Postfix NULL Terminator
- Postfix, Null Terminate, and Backslash
- Relative Path Traversal
- Client-Controlled Environment Variables
- User-Supplied Global Variables (DEBUG=1, PHP Globals, and So Forth)
- Session ID, Resource ID, and Blind Trust
- Analog In-Band Switching Signals (aka "Blue Boxing")
- Attack Pattern Fragment: Manipulating Terminal Devices
- Simple Script Injection
- Embedding Script in Nonscript Elements
- XSS in HTTP Headers
- HTTP Query Strings
- User-Controlled Filename
- Passing Local Filenames to Functions That Expect a URL
- Meta-characters in E-mail Header
- File System Function Injection, Content Based
- Client-side Injection, Buffer Overflow
- Cause Web Server Misclassification
- Alternate Encoding the Leading Ghost Characters
- Using Slashes in Alternate Encoding
- Using Escaped Slashes in Alternate Encoding
- Unicode Encoding
- UTF-8 Encoding
- URL Encoding
- Alternative IP Addresses
- Slashes and URL Encoding Combined
- Web Logs
- Overflow Binary Resource File
- Overflow Variables and Tags
- Overflow Symbolic Links
- MIME Conversion
- HTTP Cookies
- Filter Failure through Buffer Overflow
- Buffer Overflow with Environment Variables
- Buffer Overflow in an API Call
- Buffer Overflow in Local Command-Line Utilities
- Parameter Expansion
- String Format Overflow in syslog()





Attack pattern 1: Make the client invisible

- Remove the client from the communications loop and talk directly to the server
- Leverage incorrect trust model (never trust the client)
- Example: hacking browsers that lie

Washington (IAD) to Paris (CDG)				Monday, Mar 12	
Flight info	Dates	Misc	Fares		
United Airlines 914	Mar 12 5:35 pm depart IAD	stops: Non-stop	Class: Coach Fare Rules	Delete	
Boeing 777	Mar 13 7:00 am arrive CDG				
Paris (CDG) to Washington (IAD)				Tuesday, Mar 13	
Flight info	Dates	Misc	Fares		
United Airlines 915	Mar 13 1:00 pm depart CDG	stops: Non-stop	Class: Coach Fare Rules	Delete	
Boeing 777	Mar 13 3:30 pm arrive IAD				

Total Airfare (including taxes): USD 2090.76

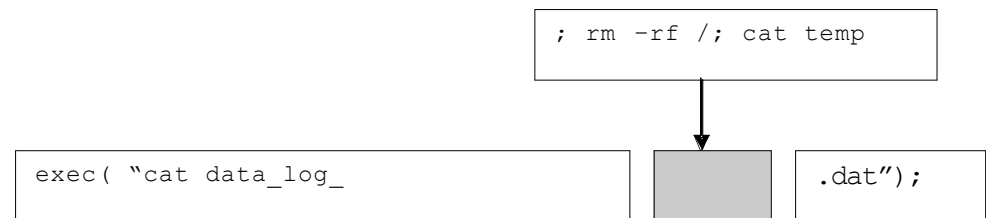
[Click to Select Your Seats](#)



Attack pattern 2: Command delimiters

- Use off-nominal characters to string together multiple commands
- Example: shell command injection with delimiters

```
<input type=hidden name=filebase  
value="bleh; [command]">
```



```
cat data_log_; rm -rf /; cat  
temp.dat
```

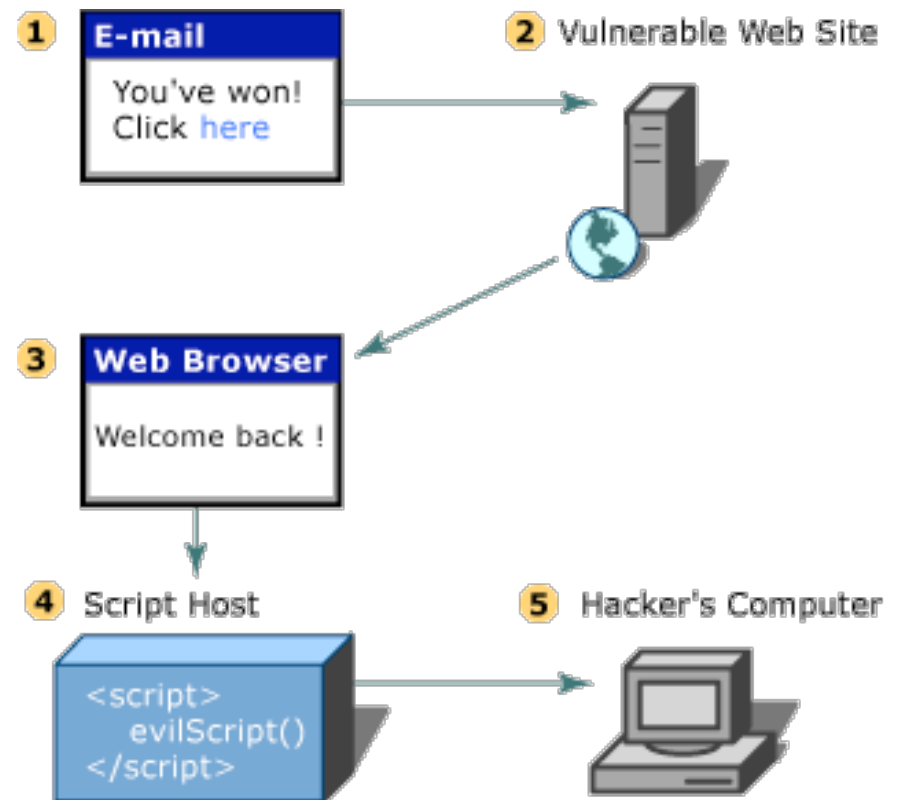
Attack pattern 3: Cross site scripting

■ XSS

- Attacker sends active content to a victim
- Content invokes a script on the vulnerable website
- Later invoked by a web browser hitting the website
- The script runs
- Attacker allowed access

■ Examples

- Javascript injection
- Inject in non-script elements
- HTTP headers
- Query strings

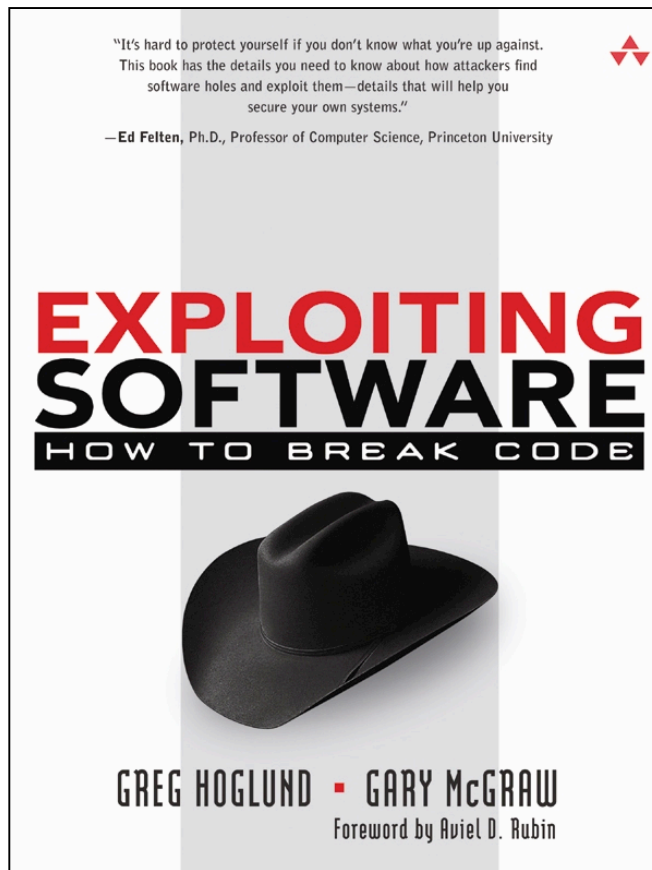




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Breaking stuff is important

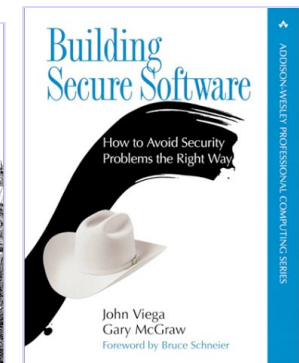
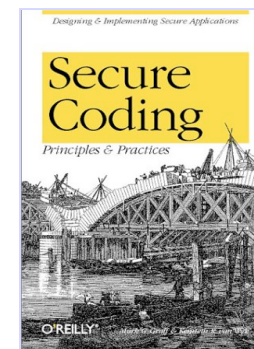


- Learning how to think like an attacker is essential
- Do not shy away from carrying out attacks on your own stuff
 - Engineers learn from stories of failure
- Attacking is fun! Fun is good!



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Software security touchpoints

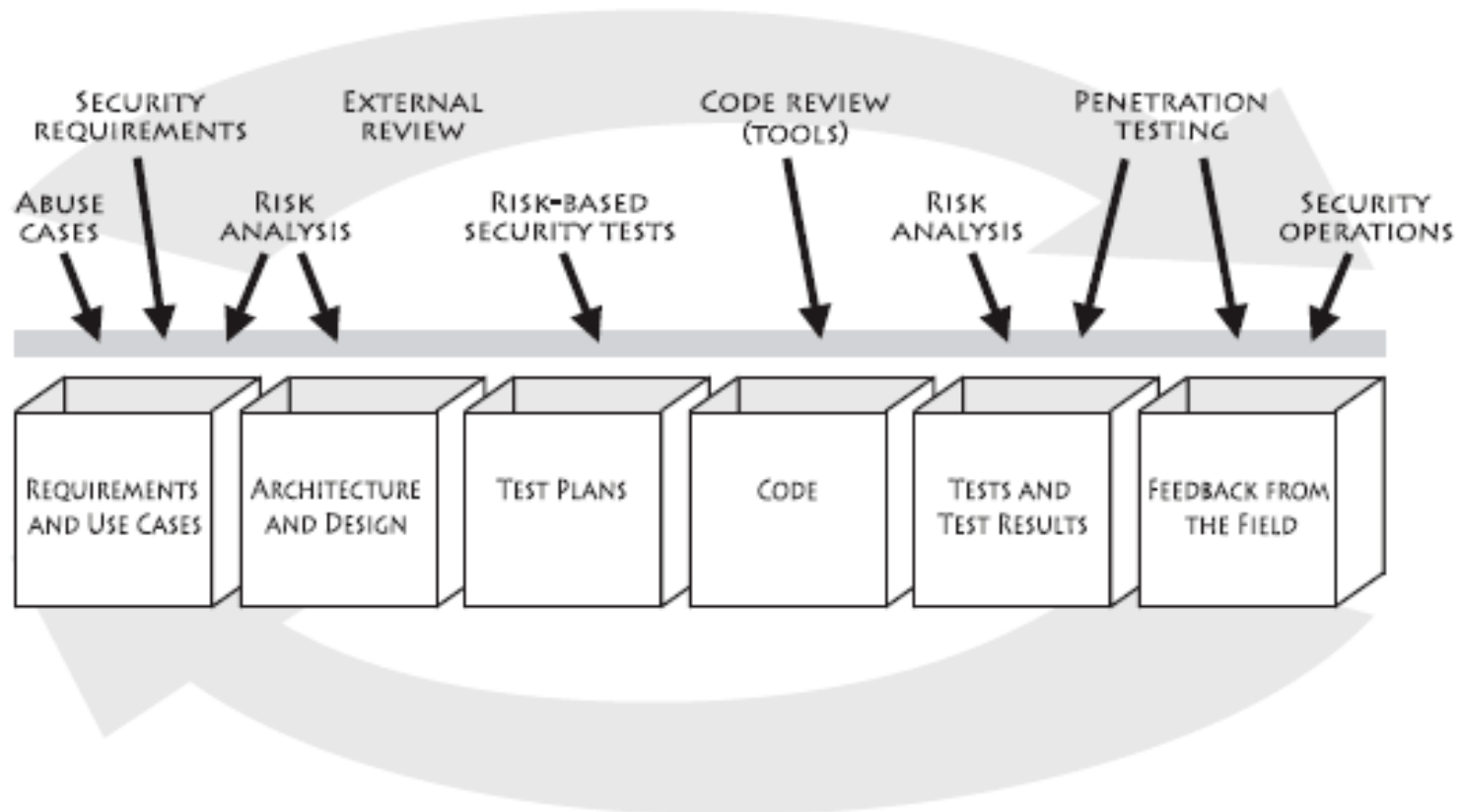


Software security touchpoints



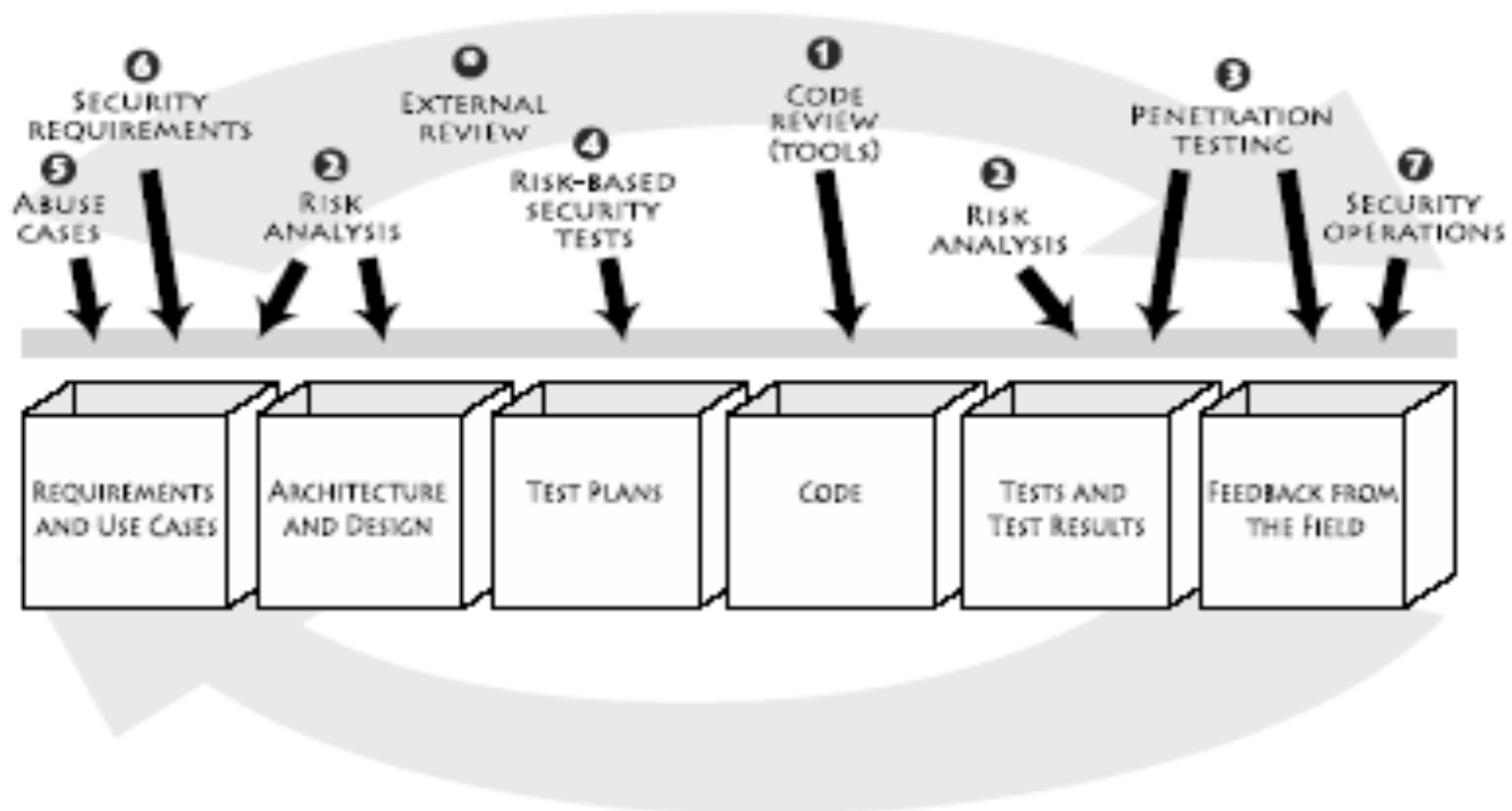


Software security touchpoints





Adopting the touchpoints

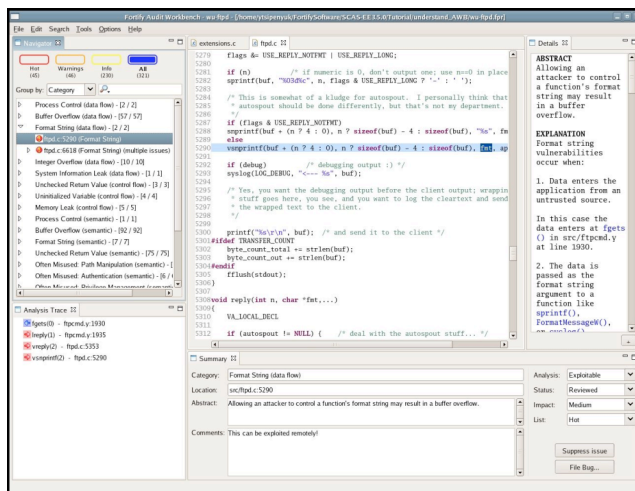




Touchpoint 1: code review (with a tool)

- Code review is a necessary evil
- Better coding practices make the job easier
- Automated tools help catch silly errors
 - Fortify/SCA (Cigital rules)
- Implementation errors do matter
 - Buffer overflows can be uncovered with static analysis
 - Static analysis
 - C/C++
 - Java
 - .NET
 - PSQL

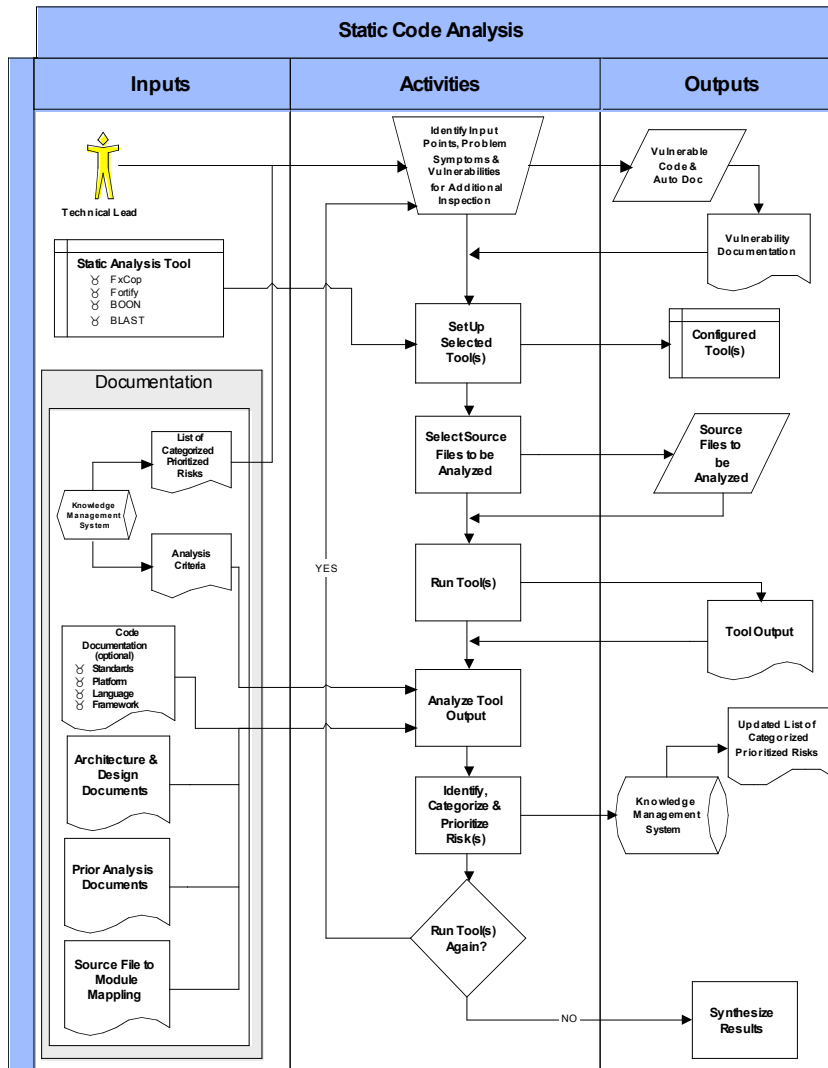
- Tracing back from vulnerable location to input is critical





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TP1: Code review

- There are many ways to apply code review technology
- Use a tool
- Integrate into the build



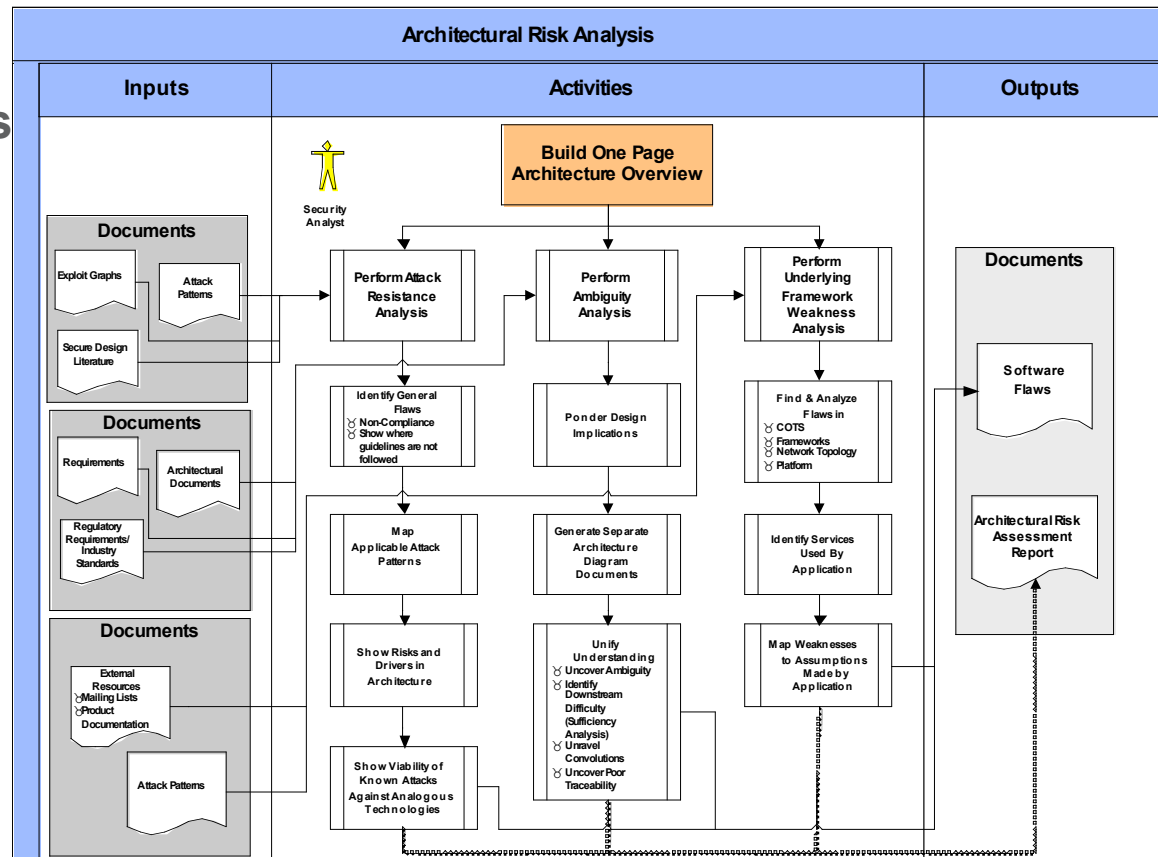


Touchpoint 2: Architectural risk analysis

- To assess and understand the risks, ask questions:
 - What is the likelihood of an attack?
 - What does the software do to support your organization's mission?
 - Is there a disaster recovery plan?
 - What would the impact be if the software were unavailable?
 - What is a tolerable down time?
- Whom should you ask?
 - Software owner
 - IT manager
 - Key users

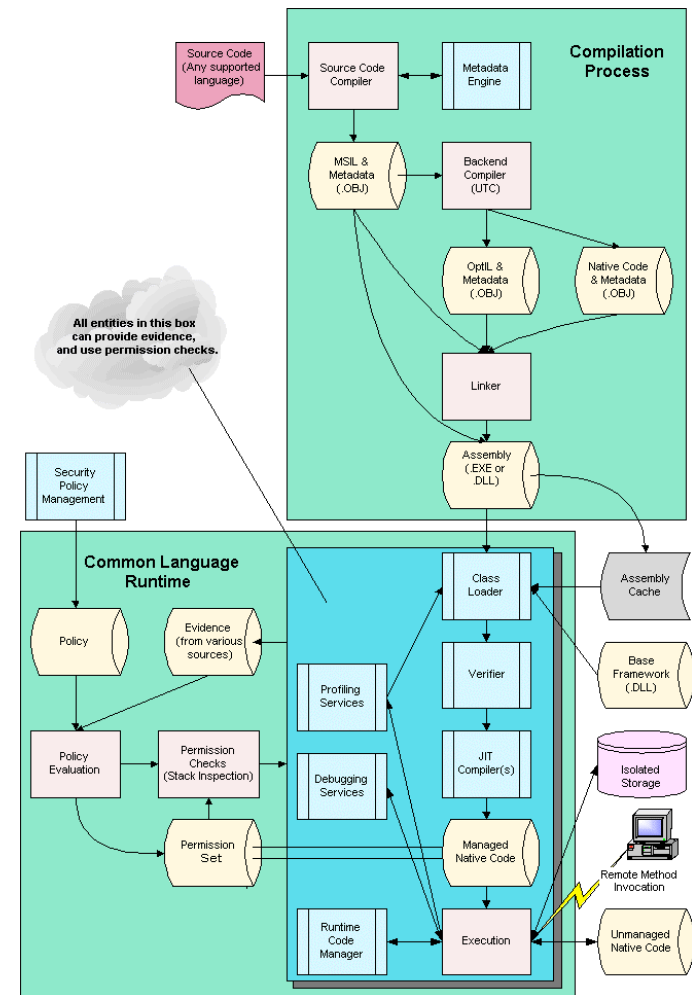
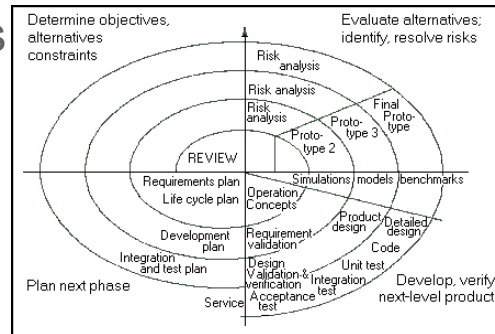
TP2: Architectural risk analysis

- Follow a process
- Build an overview (one page)
- Three steps
 - Attack resistance analysis
 - Ambiguity analysis
 - Weakness analysis
- Rank risks
- Build mitigations



TP2: Architectural risk analysis

- Designers should not do this
- Build a one page white board design model (like that →)
- Use hypothesis testing to categorize risks
 - Threat modeling/Attack patterns
- Rank risks
- Tie to business context
- Suggest fixes
- Repeat





TP2 step: Attack resistance

- Identify general flaws
 - Non-compliance
 - Where guidelines are not followed
- Map applicable attack patterns
- Identify risks in architecture
- Consider known attacks against similar technologies
- Attack Patterns
 - Pattern language
 - Database of patterns
 - Actual flaws from clients
- Exploit Graphs
 - Ease mitigation
 - Demonstrate attack paths
- Secure design

Example flaws from experience...

- Transparent authentication token generation/management
- Misuse of cryptographic primitives
- Easily subverted guard components, broken encapsulation
- Cross-language trust/privilege issues



TP2 step: Ambiguity analysis

- Consider implications of design
- Generate separate arch. diagrams
- Unify understanding
 - Uncover ambiguity
 - Identify downstream difficulty (traceability)
 - Unravel convolution
- Apprenticeship model
- Use system, technology experts
 - Win32 knowledge
 - JVM/managed code
 - Language/compiler knowledge
- Previous experience

Example flaws from experience...

- Protocol, authentication problems
- Javacard applet firewall, inner class issues, instantiation in C#
- Type safety and type confusion
- Password retrieval, fitness and strength



TP2 step: Weakness analysis

- Consider systemic flaws
 - COTS
 - Frameworks
 - Network topology
 - Platform
- Identify services
- Map weaknesses to assumptions
- Experience base
 - Assessments of COTS and platforms
- Attack patterns
- Other resources
 - Mailing lists
 - Product documentation

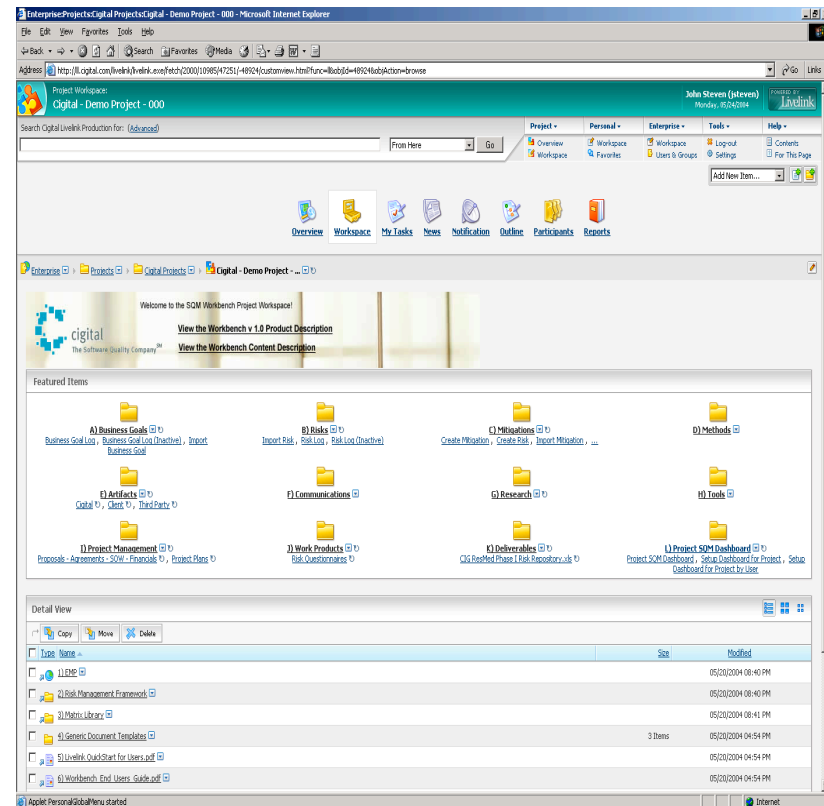
Example flaws from experience...

- Browser and other VM sandboxing failures
- Insecure service provision: RMI, COM, etc.
- Debug (or other operational) interfaces
- Unused (but privileged) product “features”
- Interposition attacks: DLLs, library paths, client spoofing



TP2: Keep track of risks

- The key to making a process like the one we described work is to **KEEP TRACK** of what you've found
- Use excel if you have nothing better
- Cigital uses the Cigital workbench
- Remember the RMF? Use it!





Touchpoint 3: Penetration testing

- A very good idea since software is bound in an environment
- How does the complete system work in practice?
 - Interaction with network security mechanisms
 - Firewalls
 - Applied cryptography
- Penetration testing should be driven by risks uncovered throughout the lifecycle
- Not a silver bullet!





Touchpoint 4: Security testing

- Test security functionality
 - Cover non-functional requirements
 - Security software probing

- Risk-based testing
 - Use architectural risk analysis results to drive scenario-based testing
 - Concentrate on what “you can’t do”
 - Think like an attacker
 - Informed red teaming



TP4: Risk-based testing

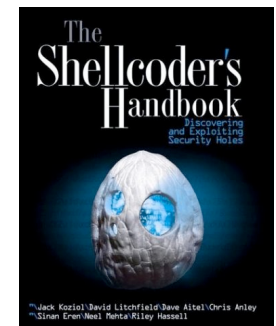
- Identify areas of potential risk in the system
 - Requirements
 - Design
 - Architecture
- Use abuse cases to drive testing according to risk
- Build attack and exploit scenarios based on identified risks
- Test risk conditions explicitly
- Example: Overly complex object-sharing system in Java Card





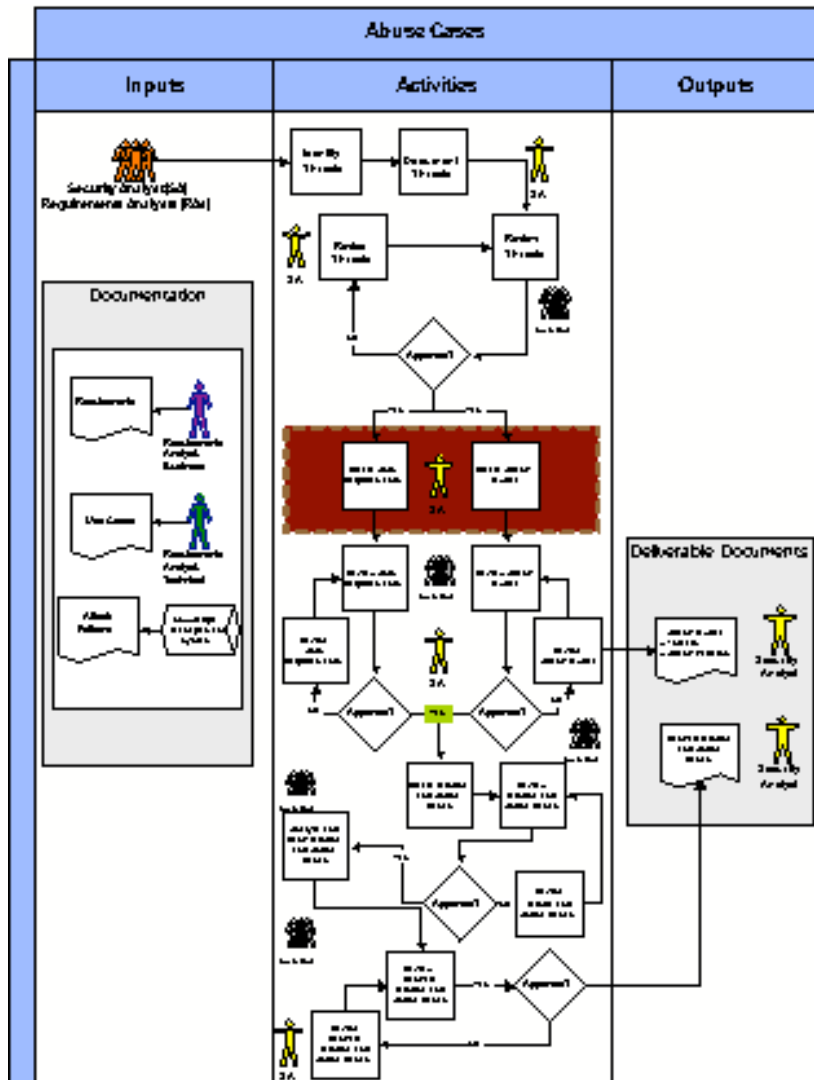
Touchpoint 5: Abuse cases

- Use cases formalize normative behavior (and assume correct usage)
- Describing non-normative behavior is a good idea
 - Prepare for abnormal behavior (attack)
 - Misuse or abuse cases do this
 - Uncover exceptional cases
- Leverage the fact that designers know more about their system than potential attackers do
- Document explicitly what the software will do in the face of illegitimate use
- Think like an attacker!

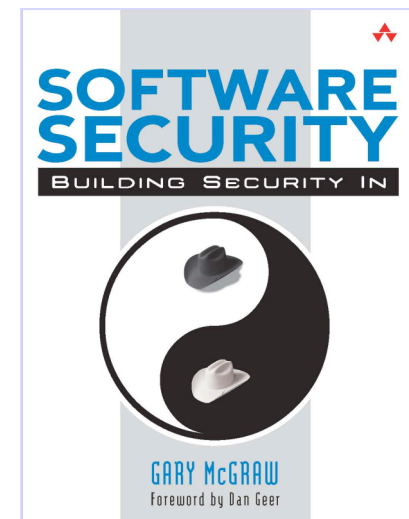




TP5: Abuse cases



- Starting with attack patterns, requirements and use cases
- Identify anti-requirements
- Build an attack model
- Determine misuse and abuse cases





Touchpoint 6: Security requirements

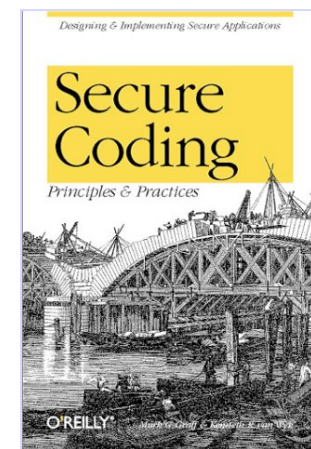
- Some security functionality maps naturally to clear requirements
 - Medical data should be cryptographically protected
 - Strongly authenticate users
 - Meet GLBA regulatory guidelines
- But do not forget that security is an emergent property of a complete system
 - An attacker needs to find only one hole
 - “Do not allow buffer overflows” is not much of a requirement!
 - “Make it secure” is vague



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Touchpoint 7: Security operations

- Use your resources!
- Network security people know an awful lot about real attacks
- Involve knowledgeable security people in as many touchpoint activities as possible
- Fine tune the deployed environment to the specific needs of your application
 - “Standard OS build” process is not enough



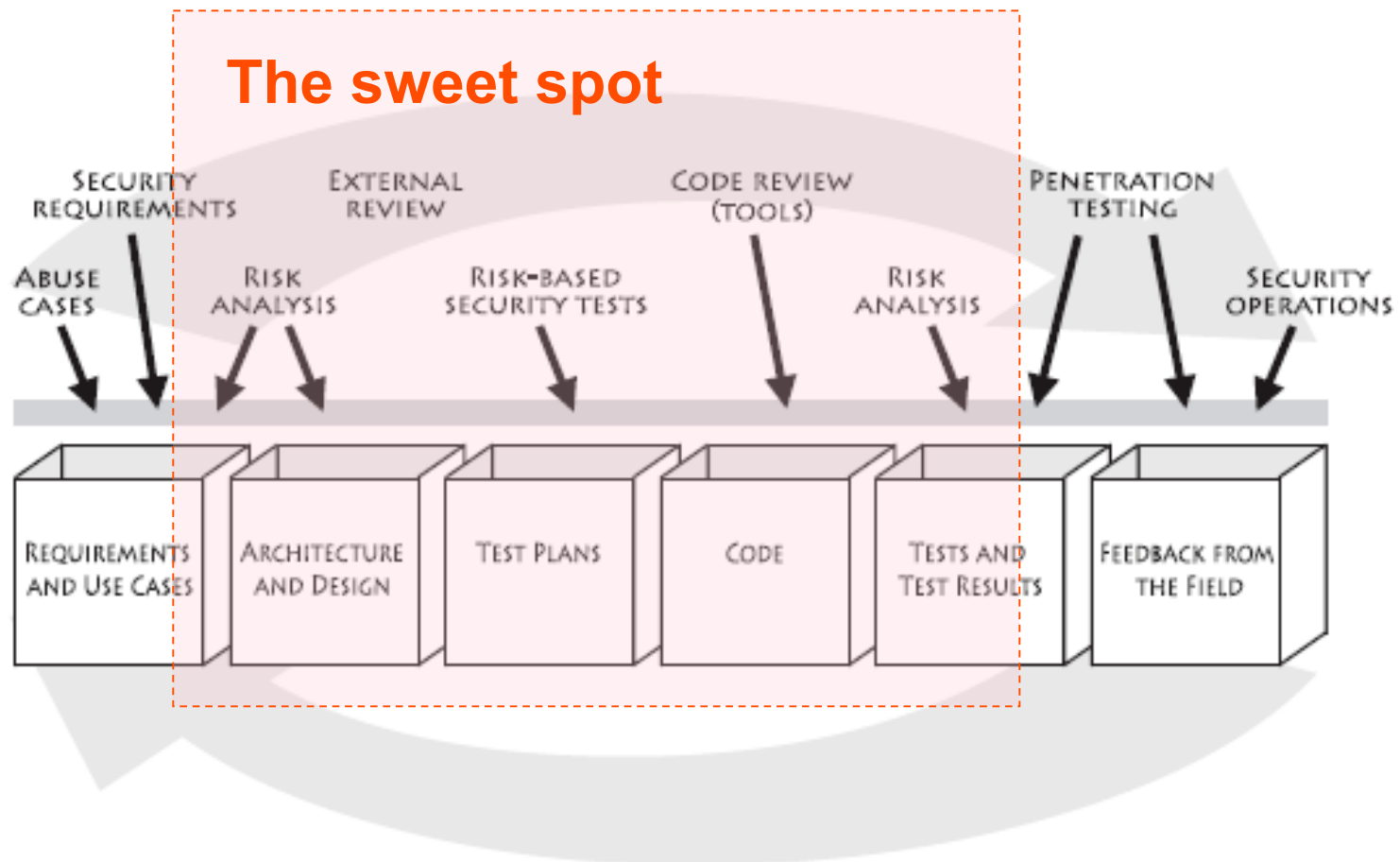


Always: External review

- Having outside eyes look at your system is essential
 - Designers and developers naturally get blinders on
 - External just means outside of the project
 - This is knowledge intensive
- Outside eyes make it easier to “assume nothing”
 - Find assumptions, make them go away
- Red teaming is a weak form of external review
 - Penetration testing is too often driven by outside → in perspective
 - External review must include architecture analysis
- Security expertise and experience really helps



Software security touchpoints



Reprise

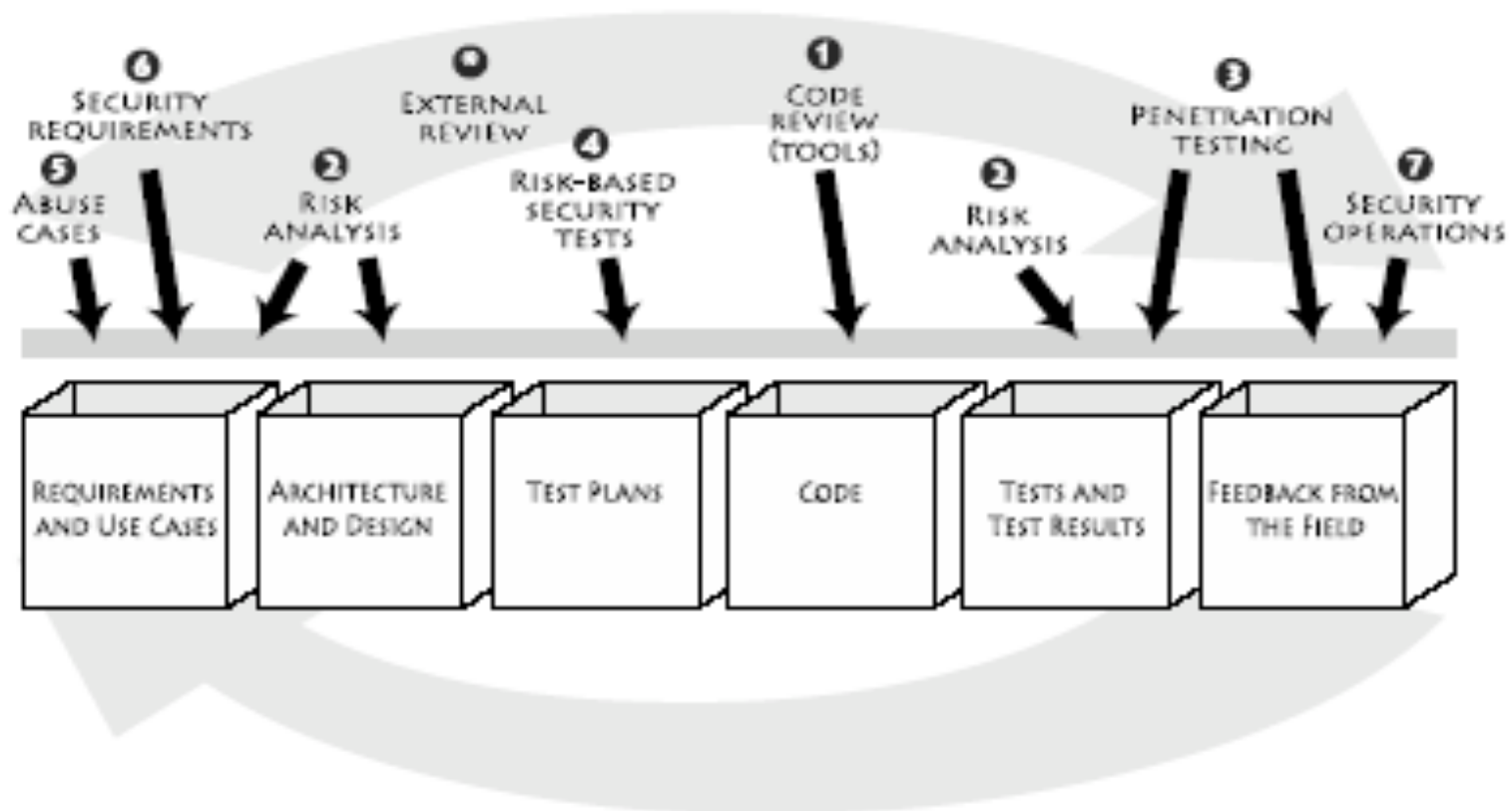


Best practices reprise

- These best practices should be applied throughout the lifecycle
 - Tendency is to “start right” (penetration testing) and declare victory
 - Not cost effective
 - Hard to fix problems
 - Start as far to the left as possible
- Abuse cases
 - Security requirements analysis
 - Architectural risk analysis
 - Risk analysis at design
 - External review
 - Test planning based on risks
 - Security testing (malicious tests)
 - Code review with static analysis tools



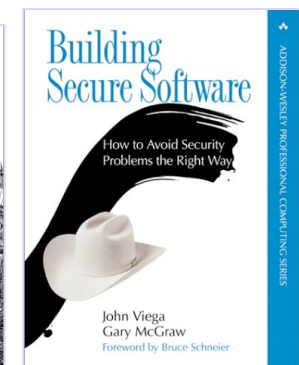
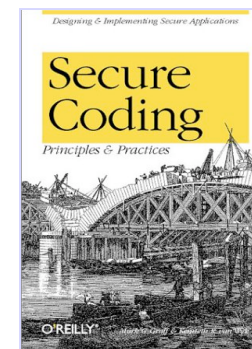
Adopting the touchpoints





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Seven pernicious kingdoms





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Outline

- Classic Pitfalls
- Seven Kingdoms
- Static Analysis and Code Review



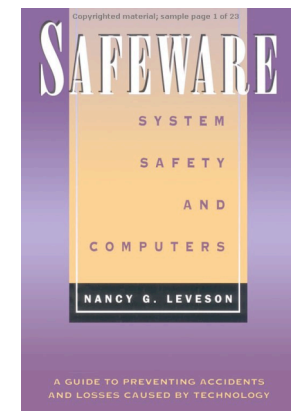
Classic Pitfalls



Learn from history

Those who cannot remember the past are condemned to repeat it.
-- Santayana

- Other engineering disciplines overcome failures by collecting failure data and analyzing failures for commonality that could lead to avoidance of that kind of failure in the future
- Failure data in software is generally considered proprietary
 - Most failure data from product development is not available for open research





Same old mistakes

- By understanding software security risks, developers can avoid them when writing their own code
- Learn by considering examples
 - Configuring applications
 - Scripts
 - Errors
 - Design flaws
- Many of the same problems crop up year after year
- Basic science to classify and categorize these problems has yet to be done
 - Bugs: implementation
 - Flaws: higher-level



Seven Kingdoms



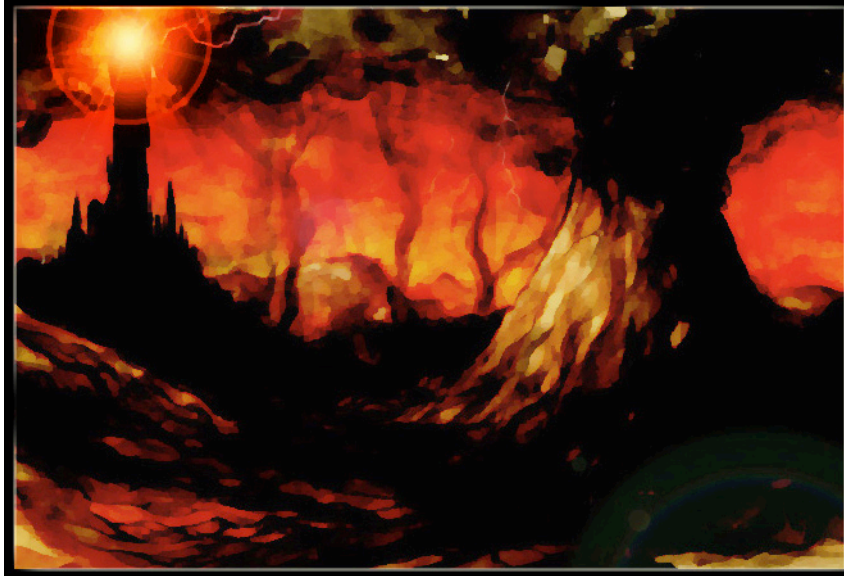
Seven pernicious kingdoms

- Input validation and representation
- API abuse
- Security features
- Time and state
- Error handling
- Code quality
- Encapsulation
- Environment



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1. Input Validation and Representation





Pernicious kingdom one

■ Input Validation and Representation

- Problems due to metacharacters, alternate encodings, numeric representations, and trusting input

- Example: **Buffer Overflow** phylum

```
int main(char ** argv, int argc) {  
    char buf[10];  
    strcpy(buf, argv[1]);  
}
```



The number one coding snafu

- “Scrubbing” user input pitfalls to avoid
 - SQL Insertion
 - Cross-Site Scripting (XSS)
 - Format string vulnerabilities
 - Integer overflows
 - Buffer overflows
 - Not a security problem per se in Java due to strict variable range enforcement
- Not a trivial issue, as complexity and subtlety abounds



Buffer overflows

- Pervasive problem, primarily in C and other non-type-safe (sometimes called “unmanaged”) code
- Responsible for huge percentage of reported vulnerabilities today
- Exploited by some of the most damaging worms
 - 1988: Morris worm
 - 2001: Code Red
 - Others: Slammer, Blaster, Sasser, Zotob



Buffer overflow causes

- String manipulation libraries
 - Flawed libc functions: strcpy, strcat, ...
 - Multibyte characters
 - Null termination errors
- Off by one errors
- Array manipulation
- Pointer arithmetic
- Others
 - Format strings
 - Integer overflow
- These all relate to reliability as well as security



Historic example: the Morris worm of 1988

- Cornell grad student Robert Tappan Morris's "Internet worm" exploited a bug in the (then) popular BSD fingerd daemon
- The vulnerable fingerd contained the following code:

```
char line[512];  
line[0] = "\\0";  
gets(line);
```

- 512 characters should be enough, shouldn't it?



Same issue in C++

- Although the `gets()` function was known to be horribly flawed for years, the same mistake was made in C++

```
char buf[BUFSIZE];  
cin >> (buf);
```

- Those cows come home yet?



Problematic function: *strcpy()*

- Although not quite as bad as *gets()*, it's darn close

```
int main(char ** argv, int argc) {  
    char buf[10];  
    strcpy(buf, argv[1]);  
}
```



Problematic function: *sprintf()*

- As with the likes of *strcpy()*, you can use *sprintf()* safely, but it isn't easy
- Is the following good or bad? (we already know it's ugly)

```
char buf[42];  
sprintf(buf, "Val1=%.8s Val2=%.8s Val3=%.8s",  
val1, val2, val3);
```



What's the deal with the n functions?

- Although the bounded versions of string functions, like `strncpy()`, are better, there's still room for silly mistakes
- Truncation can cause odd behavior
- Example: One simple mistake is to bound the data to the src buffer, as in this example from MSDN

```
int main(int argc, char *argv[]) {  
  
...  
char DirSpec[MAX_PATH + 1];  
printf("Target dir is %s.\n", argv[1]);  
strncpy(DirSpec, argv[1], strlen(argv[1])+1);
```



Problematic function: *strncat()*

- Example: The *strncat()* function is misleading because it doesn't accept a bound on the total size of the destination buffer, but rather the remaining space available in the destination buffer

```
char* buf[512];  
strcpy(buf, "The argument is");  
strncat(buf, argv[1], 512);
```



Format string vulnerabilities

- Format string vulnerabilities occur when an attacker can control a format string
- Although not technically buffer overflows, they almost invariably lead to read/writes outside a buffer's bounds
 - Including execution of arbitrary code placed on stack by the attacker
- First seen around 1999, but in its first full year resulted in many root exploits
 - Wu-ftp 2.*
 - Linux rpc.statd
 - Qualcomm qpopper 2.53
 - Apache + PHP3
 - BSD chpass
 - OpenBSD fstat



Format strings: root cause

- Misuse of formatting functions
 - A programmer wants to print a string
 - Which is correct?

```
printf("%s", string);  
printf(str);
```
- If an attacker can control the format string, then %n can be used to write arbitrary values anywhere in memory
- Exploits then work the same way as traditional buffer overflows
 - Overwrite return address
 - Function pointer
 - Other important values



Example: *wuftp* 2.6.0

- Widely publicized format string vulnerability occurs in the *vreply()* function, which looks much like this

```
while (fgets(buf, sizeof buf, f)) {  
    lreply(200, buf);  
    ...  
}
```

```
void lreply(int n, char *fmt, ...) {  
    char buf[BUFSIZ];  
    ...  
    vsnprintf(buf, sizeof buf, fmt, ap);  
    ...  
}
```



SQL insertion

- Problem can exist when Java or middle-tier code interacts with back-end SQL-based database
- User inputs must be pedantically screened for SQL code
 - White space, quotes, etc., are indicators
- Regular Expression (regex) filtering is key



Problem: SQL insertion

- Can enable attacker to execute arbitrary SQL commands on back-end database
- PHP/SQL Example:
 - PHP code inputs USERNAME and PASSWORD and passes to SQL back-end
 - USERNAME is entered as **bob**
 - PASSWORD is entered as **' or USERNAME='bob**
 - Back-end executes **Select ID from USERS where USERNAME='bob' and PASSWORD=' or USERNAME='bob'**
 - Instead of **Select ID from USERS where USERNAME='bob' and PASSWORD='password'**



SQL insertion - example

```
Pattern pattern;  
Matcher matcher;  
String regex = "\\W";
```

 [1]

```
boolean userOkay = true;  
boolean passOkay = true;
```

```
// check username field
```

```
pattern = Pattern.compile(regex);  
matcher = pattern.matcher(userField.getText());
```

 [2]

[1] Begin by defining the regular expression itself

[2] Compile the regex and apply it to the string in question

(Even better: use `PreparedStatement`)



Complications in parsing input

- Lots of things can make parsing through input fields complex
- Whitelisting and blacklisting approaches
 - Assume input is dangerous until it is proven to be safe
- Internationalization
 - Unicode can be used to obfuscate SQL insertion, XSS, etc.
 - /etc/passwd—seems easy enough to parse, right?



Unicode - example

```
Pattern pattern;  
Matcher matcher;  
String regex = "^[a-zA-Z0-9_\\-\\.]*$";  
String unicodeRegex = "[u002f]u002F";  
String file = fileField.getText();  
  
// check filename field  
  
pattern = Pattern.compile(unicodeRegex);  
matcher = pattern.matcher(file);  
  
if (matcher.find()) {  
    unicodeFound = true;  
}  
  
pattern = Pattern.compile(regex);  
matcher = pattern.matcher(file);  
  
if (matcher.find()) {  
    filenameOkay = false;  
}
```

[1]

[2]

- [1] Define a regex to search for unicode characters (u002f = “\”)
- [2] Check for specified unicode characters in the file name



Good practice: take care with config files

- Check configuration files
 - Can be ripe target for attackers
 - Verify read/write controls are safe
 - Verify data content before acting
- User inputs
 - Command line parameters and desktop icons
 - URLs
 - Assume it to be harmful until proven otherwise
- Consider also where other user inputs can come from
 - Signals, registry keys, mouse actions, and so on...



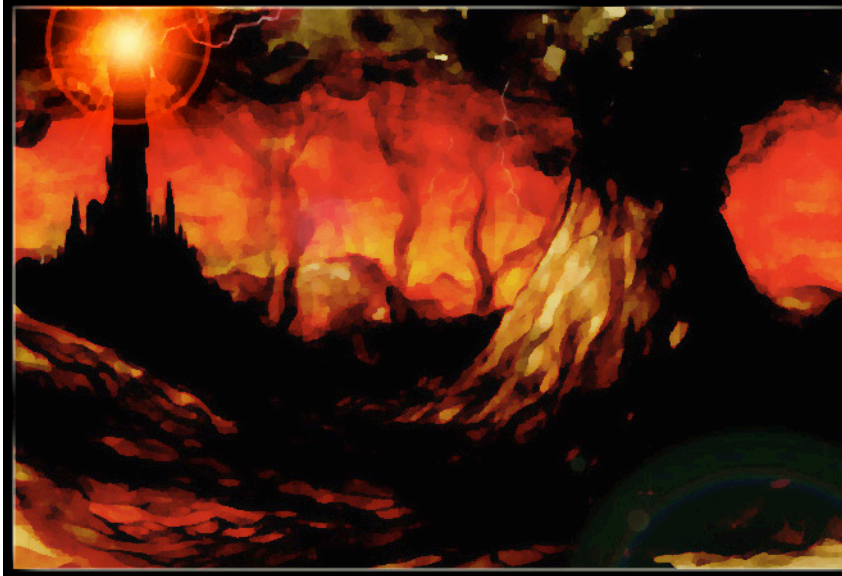
Phyla: Input validation and representation

- Buffer Overflow
- Command Injection
- Cross-Site Scripting
- Format String
- HTTP Response Splitting
- Illegal Pointer Value
- Integer Overflow
- Log Forging
- Path Traversal
- Process Control
- Resource Injection
- Setting Manipulation
- SQL Injection
- String Termination Error
- Struts: Duplicate Validation Forms
- Struts: Erroneous validate() Method
- Struts: Form Bean Does Not Extend Validation Class
- Struts: Form Field without Validator
- Struts: Plug-in Framework Not in Use
- Struts: Unused Validation Form
- Struts: Unvalidated Action Form
- Struts: Validator Turned Off
- Struts: Validator without Form Field
- Unsafe JNI
- Unsafe Reflection
- XML Validation



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2. API Abuse





Pernicious kingdom two

■ API Abuse

- A caller fails to honor the contract between the caller and the callee
- Dangerous function, Unchecked Return Value, and others
- Example: **Often Misused: Authentication** phylum

```
String ip = request.getRemoteAddr();
InetAddress addr = InetAddress.getByName(ip);
if (addr.getCanonicalHostName().endsWith(
    "trustme.com")) {
    trusted = true;
}
// Relying on DNS lookup for recognizing trusted hosts
```




Comparing Java classes

- Never make a decision based on the name of a class
 - A program may treat two classes the same when they actually differ
 - Class names are trivial to forge or substitute
 - At the very least, verify that the name being checked is within the current classloader



Example: *readlink()*

- Abuse of *readlink()*, which although it fills a string buffer does not null terminate the buffer

```
readlink(path, buf, MAXPATH);  
int length = strlen(buf);
```

- The value returned from *strlen()* is likely to be incorrect – perhaps wildly so – and may even result in a buffer overflow or other runtime erratic behavior



Example: SYN flood

- Attacker initiates, but does not complete TCP session opening protocol
- Victim's TCP stack is left in a wait state
- Attacker repeats until victim's resource pool is saturated
- Victim is now effectively off the net – DoS
- Why would someone want to do this?



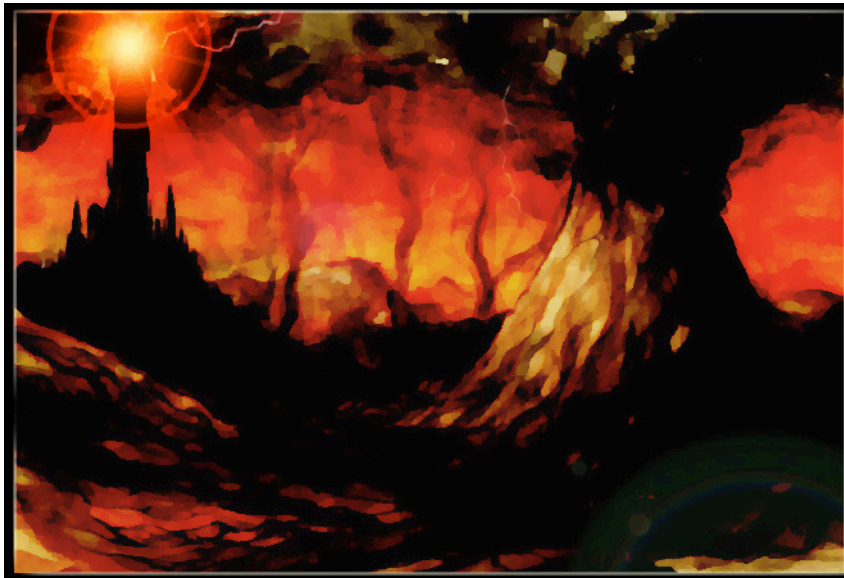
Phyla: API abuse

- Dangerous Function
- Directory Restriction
- Heap Inspection
- J2EE Bad Practices: getConnection()
- J2EE Bad Practices: Sockets
- Often Misused: Authentication
- Often Misused: Exception Handling
- Often Misused: Path Manipulation
- Often Misused: Privilege Management
- Often Misused: String Manipulation
- Unchecked Return Value



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3. Security Features





Pernicious kingdom three

■ Security Features

- Poorly handled authentication, access control, confidentiality, cryptography, and privilege management
- Insecure Randomness, Password Management, Privacy Violation, and others
- Example: **Privacy Violation** phylum

```
id = getId();  
pass = getPassword();  
type = getType();  
tstamp = getTimestamp();  
... // Private info leaking into a log file  
dbmsLog.log(id+":"+pass+":"+type+": "+tstamp);
```



Signing JAR files

- Signing JAR files can be dangerous
- A signed JAR might be trusted more than is warranted
 - But, signed JARs also can be useful
 - Authentication and integrity checking
- If you must sign, put your signed classes into one JAR file, all by themselves



Storing secrets

- “Hard coding” sensitive information in source code is dangerous
 - Class file can be viewed
 - Class de-compilers (e.g., jode) can expose



Storing secrets - example

```
public class SecretInfo {  
    public static void main(String[] args) {  
        String username = "bassethoundsruletheworld";  
        String password = "t0ps3cr3t!";  
        /* do something with username and password */  
    }  
}
```

[1]

[1] Two strings are defined that contain sensitive information



Storing secrets – example (cont'd)

```
/* SecretInfo - Decompiled by JODE
 * Visit http://jode.sourceforge.net/
 */
public class SecretInfo
{
    public static void main(String[] strings) {
        String string = "bassethoundsruletheworld";
        String string_0_ = "t0ps3cr3t!";
    }
}
```

[2]

[2] Using a decompiler, the values of both strings can be retrieved from a compiled .class file



Privilege handling

- Don't forget the principle of least privilege
 - Avoid privileged code if at all possible
- Tips
 - Design things so that program does not need privileges
 - Develop code without privileges enabled
- Did you know?
 - 90% of Windows software can't be installed without Administrator privileges
 - 70% can't be run without Administrator privileges
 - 10,000 lemmings can't be wrong!



Why privileges are needed

- Interact directly with hardware
- Other shared resources
 - Network ports, config, registry
- Alter OS behavior
- Override file system protections
 - Install new files
 - Update protected files
 - Access files that belong to other users



Case study: *lpr*

- Redhat lpr (Oct 1999)
- Setuid root in order to talk to printer device

```
int fd;
for (int i=1; i < argc; i++) {
    /* first make sure that the user can read the
    file, then open it */
    if (!access(argv[i], O_RDONLY)) {
        fd = open(argv[i], O_RDONLY);
    }
    print(fd);
}
```



Case study: *lpr*

- File access race condition! Fix:

```
int fd;
for (int i=1; i < argc; i++) {
    int uid = getuid(); int gid = getgid();
    int original_euid = geteuid();
    int original_egid = getegid();
    seteuid(uid); setegid(gid);
    fd = open(argv[i], O_RDONLY);
    seteuid(original_euid);
    setegid(original_egid);
}
print (fd);
```



Case study: lpr

- Do you think that it's fixed now?
- No! *seteuid()* return value ignored
- No one expects *seteuid()* to fail since we're root

- POSIX capabilities vulnerability (June 2000)
- Attackers can cause *seteuid()* call to fail

- Not so simple, is it?



When are random numbers needed?

- Some numbers need to be cryptographically secure
 - Crypto applications
 - Generated passwords
 - Port randomization
 - External unique identifiers such as session tokens
 - Discount codes
- Some do not
 - Monte Carlo simulation systems
 - Internal unique identifiers



Example: Security depends on unpredictability

- The following code generates “unique” identifiers for online users who make a purchase. Because *lrand48()* is a statistical PRNG, it is easy for an attacker to predict

```
char* CreateReceiptURL() {
    int num; time_t t1;
    char *URL = (char*) malloc(MAX_URL);
    if (URL) {
        (void) time(&t1);
        srand48((long) t1)
        sprintf(URL, "%s%d%s, http://test.com,
                lrand48(), ".html");
    }
    return URL; }
```



Choosing a PRNG

- Hardware can be good, if available
- OS may provide good random sources
 - /dev/urandom is almost always the right choice for user apps
 - /dev/random blocks and may be exhausted since shared
- Current state of the art
 - Fortuna (described in Schneier's Practical Cryptography)
 - Implementations
 - Win C++ (<http://www.citadelsoftware.ca/fortuna/Fortuna.htm>)
 - Linux /dev/urandom driver (<http://jlcooke.ca/random>)
- Freebie in Microsoft-friendly code
 - CryptoGenRandom()



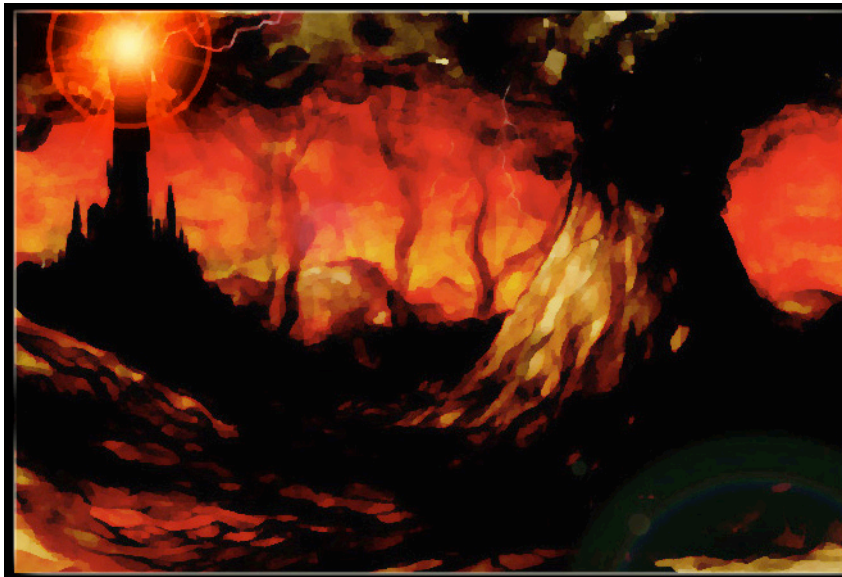
Phyla: Security features

- Insecure Randomness
- Least Privilege Violation
- Missing Access Control
- Password Management
- Password Management: Empty Password in Configuration File
- Password Management: Hard-Coded Password
- Password Management: Password in Configuration File
- Password Management: Weak Cryptography
- Privacy Violation



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4. Time and State





Pernicious kingdom four

■ Time and State

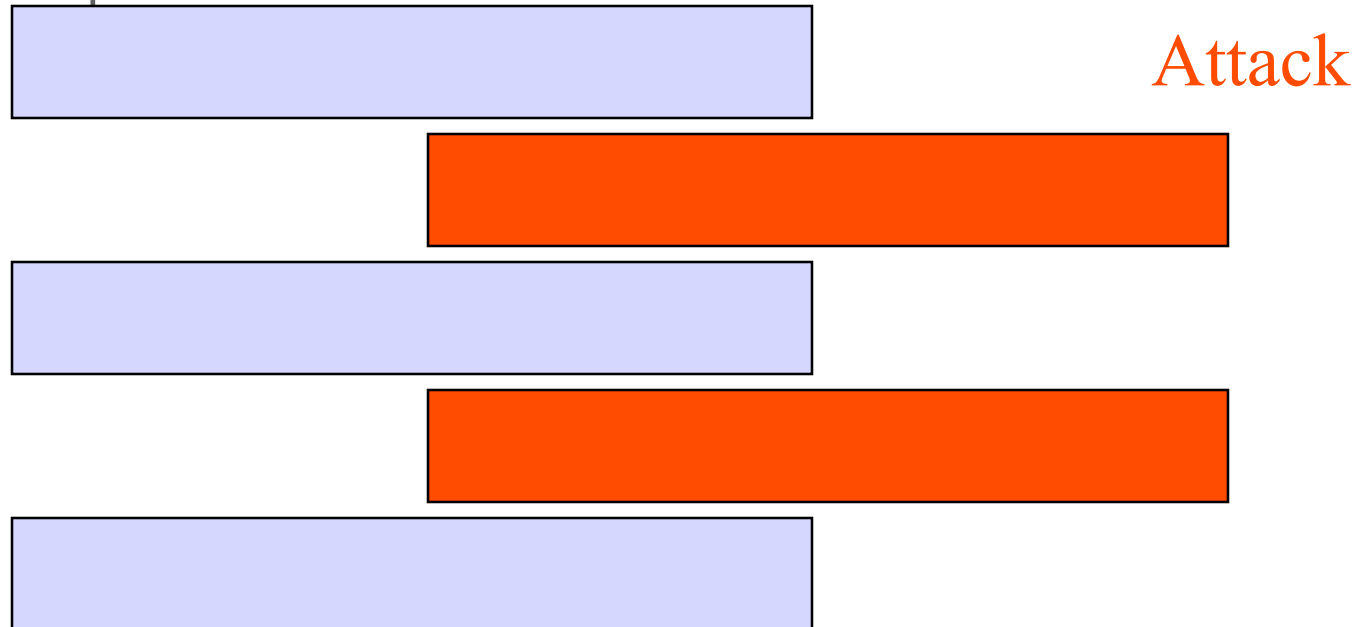
- Unexpected interactions between threads, processes, time, and information that happen through shared state: semaphores, variables, file system, etc.
- File Access Race Condition TOCTOU, Deadlock, and others
- Example: **Session Fixation** phylum

```
private void auth(LoginContext lc,  
HttpSession session)  
    throws LoginException {  
    ... // No call to session.invalidate()  
    lc.login();  
    ...  
}
```



RISK: Race condition

- Time makes all the difference
- Atomic operations that are not atomic





A simple (broken) Java servlet

```
import java.io.*;
import java.servlet.*;
import java.servlet.http.*;
public class Counter extends HttpServlet{
    int count = 0;
    public void doGet(HttpServletRequest in, HttpServletResponse out)
        throws ServletException, IOException {
        out.setContentType("text/plain");
        PrintWriter p = out.getWriter();
        count++;
        p.println(count + " hits so far!");
    }
}
```

Race condition



A simple (fixed) Java servlet

```
import java.io.*;
import java.servlet.*;
import java.servlet.http.*;
public class Counter extends HttpServlet{
    int count = 0;
    public synchronized void
        doGet(HttpServletRequest in, HttpServletResponse out)
            throws ServletException, IOException {
        out.setContentType("text/plain");
        PrintWriter p = out.getWriter();
        count++;
        p.println(count + " hits so far!");
    }
}
```




TOCTOU

- Race conditions on Unix files are famous
- Passwd example
 - Step 1: open file and read it in
 - Step 2: create and open “ptmp” in same directory
 - Step 3: open password file again, copying unchanged contents into ptmp while updating
 - Step 4: Close both password file and ptmp, then name ptmp the password file
- If an attacker makes use of unix’s linking facility, an attack is possible
- Change the system state in a subtle way in order to cause the system to do something dangerous



Threads (J2EE)

- Thread management in a web application is prohibited by the J2EE standard
- Difficult and likely to produce unpredictable results such as deadlocks, race conditions and other synchronization errors
- Rather than managing threads directly, use standards such as message driven beans and EJB timer service provided by the container



Good practice: watch out for web content

- Web data
 - Watch out for data in hidden fields
 - Even though it is within page, user can still alter
- Web cookies
 - Can also be manipulated by user
 - Classic example: changing customer ID or shopping cart price totals
- State data must be protected
 - Encryption is commonly used
 - Verify that no data has been tampered with



Serialization

- Largely fixed in latest JDK versions
 - Previous default allowed serialization
 - New default requires class to implement Serializable interface
- When serialized, an object is written to disk directly, including internal memory
- If you must make something serializable, declare private data transient



Serialization - example

```
public class SerializableClass implements Serializable {  
  
    private transient String secret = "Cats are smarter than dogs.";  
    private String notSecret = "Dogs drool everywhere.";  
  
    public SerializableClass() { }  
}
```

 [1]

```
notSecretLjava/lang/String;xptDogs drool everywhere.x
```

 [2]

- [1] A serializable class that defines two private strings
- [2] Output of the serialized class when read by a simple text editor (note that the transient string is not displayed)



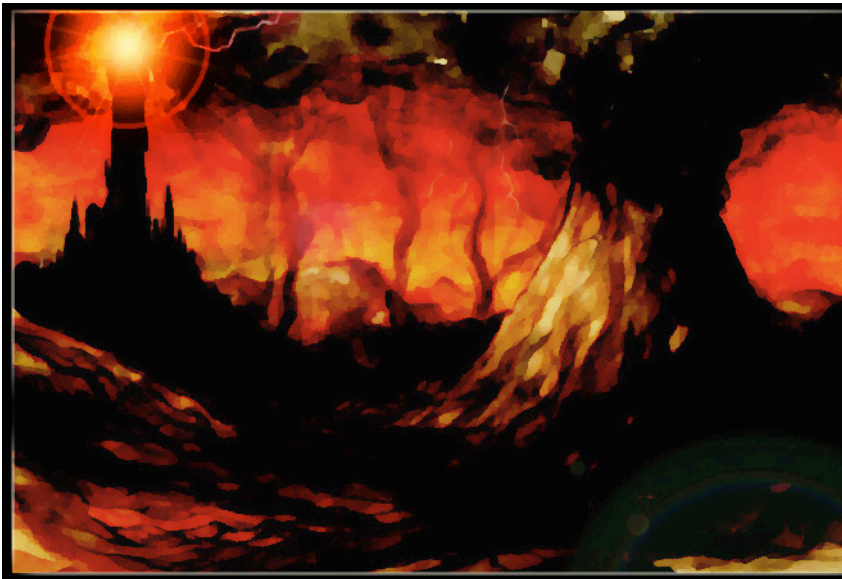
Phyla: Time and state

- Deadlock
- Failure to Begin a New Session upon Authentication
- File Access Race Condition: TOCTOU
- Insecure Temporary File
- J2EE Bad Practices: System.exit()
- J2EE Bad Practices: Threads
- Signal Handling Race Conditions



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5. Error Handling





Pernicious kingdom five

■ Error handling

- Both poor error handling and generation of errors that either leak information or are difficult to handle
- Empty Catch Block, Overly-Broad Catch Block, and others
- Example: **Empty Catch Block** phylum

```
try {  
    attemptToDoSomethingImportant();  
}  
catch (ImportantException e) {  
    // How should this exception be handled?  
}
```




Error handling: the problem

- Ignoring exceptional conditions and their ramifications
 - A symptom: failure to think about what could go wrong
 - An outcome: leads to inconsistent and unexpected program state
- Unchecked return values
- Exception handling
- Signal handling



Legacy problems

- Many standards to choose from
 - *fork()* – 0 == success
 - *strtol()* – 0 == failure
 - *strcmp()* – 0 == true
 - *issetugid()* – 0 == false
 - *fork()* -- >0 == success
- And this doesn't even address multithreaded apps
- Always check those reference manuals before assuming!



Allocation problems

- Failure to check for memory allocation failure

```
buf = (char*) malloc(req_size);
strncpy(buf, xfer, req_size);
```
- What could go wrong?
- Bad for at least three reasons
 - No opportunity to recover
 - Impossible to exit gracefully
 - No opportunity of collecting diagnostic information



Missing error handling (J2EE)

- Un-handled exceptions can provide an attacker with potentially dangerous information, such as an SQL query string, the type of database being used, or application version numbers
- Web applications should always specify default error pages and handle standard HTTP error codes



Missing error handling - example

Include the following entries in the *web.xml* file to specify default error pages...

```
<error-page>
  <exception-type>java.lang.Throwable</exception-type>
  <location>/error.jsp</location>
</error-page>
<error-page>
  <error-code>404</error-code>
  <location>/error.jsp</location>
</error-page>
<error-page>
  <error-code>500</error-code>
  <location>/error.jsp</location>
</error-page>
```



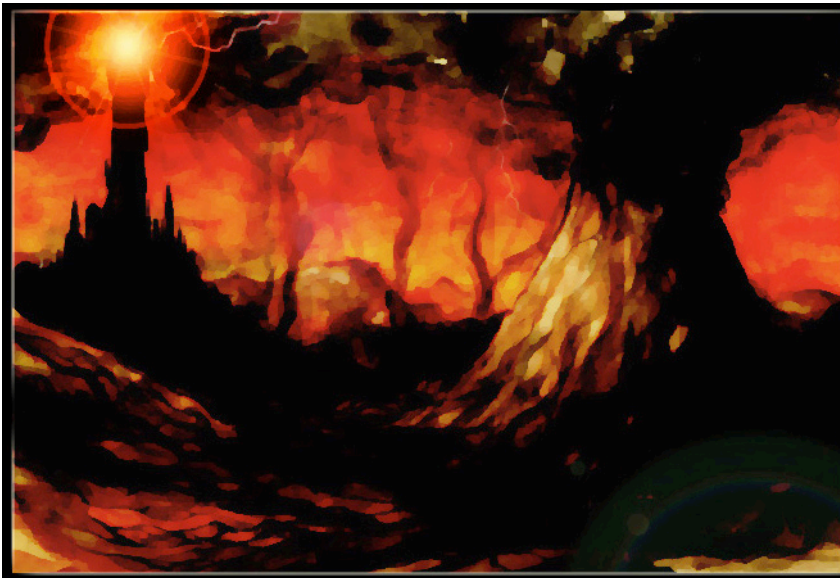
Phyla: Error handling

- Catch NullPointerException
- Empty Catch Block
- Overly Broad Catch Block
- Overly Broad Throws Declaration
- Unchecked Return Value



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6. Code Quality





Pernicious kingdom six

■ Code Quality

- Poor code quality indicates security problems likely
- Memory Leak, Null Dereference, Uninitialized Variable, and others
- Example: **Attribute Stored in HttpSession Might Not Be Serializable** phylum

```
public class MyAttribute {
    ...
}

public void add (HttpSession s, MyAttribute a) {
    session.setAttribute("attribute", a);
}
```




Code quality issues

- All have the potential to allow denial of service attacks
- More often leads to unpredictable behavior
 - Exceedingly difficult to test for
 - Read “The Bug” by Ellen Ullman
- Unpredictable behavior is the friend of the attacker



Example: memory leak

- Find easy cases with tools like Purify
- Hard cases can be dynamic flow driven and really tough to find
- Common causes: error conditions, confusion over responsibility

```
char* getBlock(int fd) {
    char* buf = (char*( malloc(BLOCK_SIZE);
    if (!buf) {
        return NULL;
    }
    if (read(fd, buf, BLOCK_SIZE) != BLOCK_SIZE) {
        return NULL;
    }
}
return buf;
```



Example: use after free

```
char* ptr = (char*) malloc(SIZM);  
...  
if (err) {  
    abrt = 1;  
    free(ptr);  
}  
...  
if (abrt) {  
    logError("operation aborted before commit", ptr);  
}
```

- And sometimes it works!
- Memory may be re-allocated by the time the error is logged



Example: double free

- Most often causes a crash, but can result in buffer overflow under rare circumstances

```
char* ptr = (char*) malloc(SIZM);
```

```
...
```

```
if (abrt) [  
    free(ptr);  
]
```

```
...
```

```
free(ptr);
```



Portability problems

- Internal buffer overflows in some implementations of *getopt()*
 - Avoid with good input validation
- In many cases, you cannot avoid problems
- Examples
 - *vfork()* behavior varies by platform
 - *strcmpi()* is not defined on many UNIX systems
 - *memmem()* problematic due to changes between versions whereby order of the arguments is reversed



Returning mutable objects

- Mutable objects are references to specific locations in memory
 - The most common example is an array
- Returning a mutable object to malicious code enables an attacker to modify the contents of memory pointed to by the object



Returning mutable objects - example

```
public static void main(String[] args) {  
    EmployeeInfo info = new EmployeeInfo("Steve", "Dallas", "Red");  
    // modify the employee name (because we can)  
    String[] localname = info.getName(); [1]  
    // Here's where the danger is; the employee's name is being  
    // stored into an array. Since that is mutable, it can be changed  
    // by an attacker, even if it is declared private.  
    localname[0] = "L33t"; [2]  
    localname[1] = "H4x0r";  
    String[] name = info.getName();  
    System.out.println("\nEmployee Name: " + name[0] + " " +  
        name[1]);  
    // will output "Employee Name: L33t H4x0r"
```

[1] Store a reference to a mutable array in a local context

[2] Modify the original array by changing the local array



Storing mutable objects

- In a similar way as returning mutable objects, storing mutable objects passed to your code can lead to problems
 - Especially if you act on the returned object(s)
- See example—MutableStorage



Public static final mutable objects

- Public static final mutable objects can still be modified, because only the reference to the object is constant



Java Initialization

- Java is supposed to initialize new variables cleanly, but it's still good practice to do so manually
 - Apart from anything else, this is just a good housekeeping



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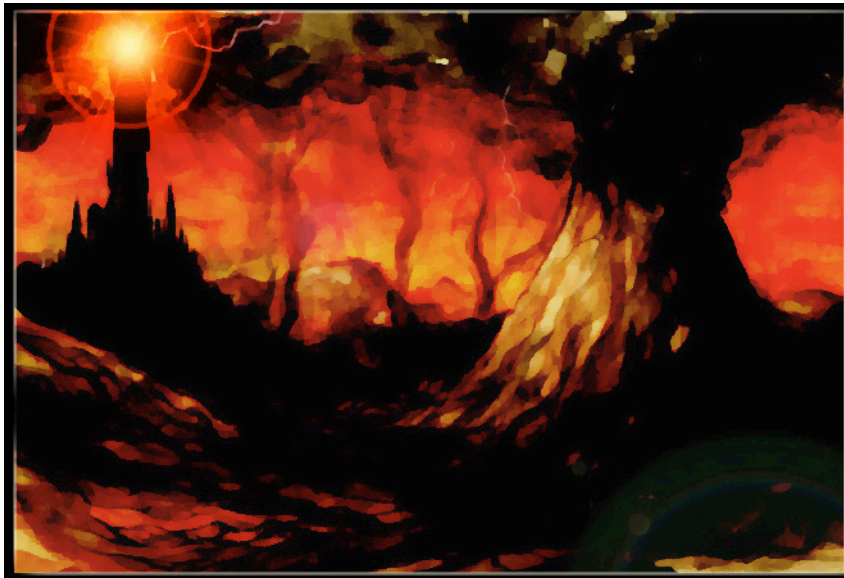
Phyla: Code quality

- Double Free
- Inconsistent Implementations
- Memory Leak
- Null Dereference
- Obsolete
- Undefined Behavior
- Uninitialized Variable
- Unreleased Resource
- Use After Free



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7. Encapsulation





Pernicious kingdom seven

■ Encapsulation

- Violation of boundaries between software components with various trust level
- System Information Leak, Trust Boundary Violation, Mobile Code: Non-Final Public Field, and others
- Example: **Field Assignment in a Servlet** phylum

```
MyServlet extends HttpServlet {                                     // Shared field
    private User user = new User();
    ...
    void getInfo(HttpServletRequest req) {
        Session s = req.getSession();
        user.userId = s.getAttribute("id");
    }
}
```



Public fields

- Public fields can be accessed by all classes
- Declare private and provide get/set methods unless they must be public
- If you absolute have to use a public field, be sure to make it final



Public fields - example

Not a good idea...

```
public class BadUserInfo {  
    public String username;  
    public String favColor;  
    public BadUserInfo(String username, String favColor) {  
        this.username = username;  
        this.favColor = favColor;  
    }  
}
```

A better idea...

```
public class GoodUserInfo {  
    private String username;  
    private String favColor;  
    public GoodUserInfo(String username, String favColor) {  
        this.username = username;  
        this.favColor = favColor;  
    }  
    public String getUsername() {  
        return username;  
    }  
}
```



Public methods

- Similarly, make sure that your methods are explicitly made private
- Prevents interface from being maliciously accessed
 - E.g., providing tainted data
- If a method must be made public, be sure to document the reason
- See example – MethodAccess



Public methods - example

```
private String username;  
private String favColor;  
  
public BadUserInfo(String username, String favColor) {  
    setUsername(username);  
  
    this.favColor = favColor;  
}  
  
public void setUsername(String user) { [1]  
    username = user;  
}
```

[1] Be sure that methods are made private unless they must be public, otherwise they can be invoked by any class



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Public static modifier

- Public static fields and methods can be accessed by other classes even if they don't instantiate



Public static modifier - example

```
public class Widget {  
    public static int height = 2; [1]  
    public int width = 10;  
  
    public Widget() {}  
}
```

```
public class StaticModifier {  
  
    public static void main(String[] args) {  
        System.out.println("\nWidget height: " + Widget.height + "\n"); [2]  
    }  
}
```

[1] The *Widget.height* field is defined as public static

[2] Any class is now able to access/modify the *height* field without instantiating the *Widget* class



Package scope

- Any class within a package can access the public and protected variables within other classes in the same package
- Thus, if you don't want to provide access to something, make it private explicitly



Package scope - example

```
package somepackage;  
  
public class Widget {  
    protected int height; [1]  
    int width;  
  
    public Widget() {  
        height = 2;  
        width = 10;  
    }  
}
```

[1] The *height* field is accessible to any class that declares itself part of the *somepackage* package



Inner classes

- The manner in which JVMs compile inner classes opens up a loophole that enables an attacker to access private members of the outer class
- Entails making creative use of the Reflection API
- See example – InnerClasses



Inner classes - example

```
public class InnerClassExample {  
    private int outerValue = 2; [1] // note the "private" modifier  
    // an inner class...  
    private class Inner {  
        private int innerValue;  
        Inner() {  
            // accessing a private field in the outer class  
            innerValue = outerValue; [2]  
        }  
    }  
}
```

[1] A private integer field is defined in the outer class

[2] The inner class accesses the private field in the outer class
(the Java compiler must create a loophole to allow this)



Inner classes – example (cont'd)

```
// call the access$000 method using reflection  
Method access = insecure.getDeclaredMethod("access$000", new Class[]  
    {insecure});  
Object value = access.invoke(null, new Object[] {example});
```

[3]

[3] The Java compiler creates a method called *access\$000* that can be called using Reflection to obtain the value of the private field



Finalization

- If methods and classes aren't made final, they can be extended in unforeseen ways and may enable an attacker to access or alter otherwise protected objects and information



Finalization - example

```
public final class Widget { [1]
    private int height;
    private int width;

    public Widget(int height, int width) {
        this.height = height;
        this.width = width;
    }

    public final int getHeight() {
        return height;
    }
}
```

- [1] Define classes to be final whenever possible to prevent them from being extended in unforeseen ways



Cloning

- If an object can be cloned, an attacker may be able to bypass its constructor, which could lead to disclosing uninitialized memory space
- If an object must implement the Cloneable interface, make sure to provide an explicit final clone() method as early in the inheritance hierarchy as possible



Cloning - example

```
class Widget {  
  
    private int height = 2;  
    private int width = 10;  
  
    public Widget() { }  
  
    public int getWidth() {  
        return width;  
    }  
  
    public final Object clone() throws java.lang.CloneNotSupportedException {  
        throw new java.lang.CloneNotSupportedException();  
    }  
}
```

[1]

[1] To prevent cloning, override the *clone()* method and throw a *java.lang.CloneNotSupportedException*



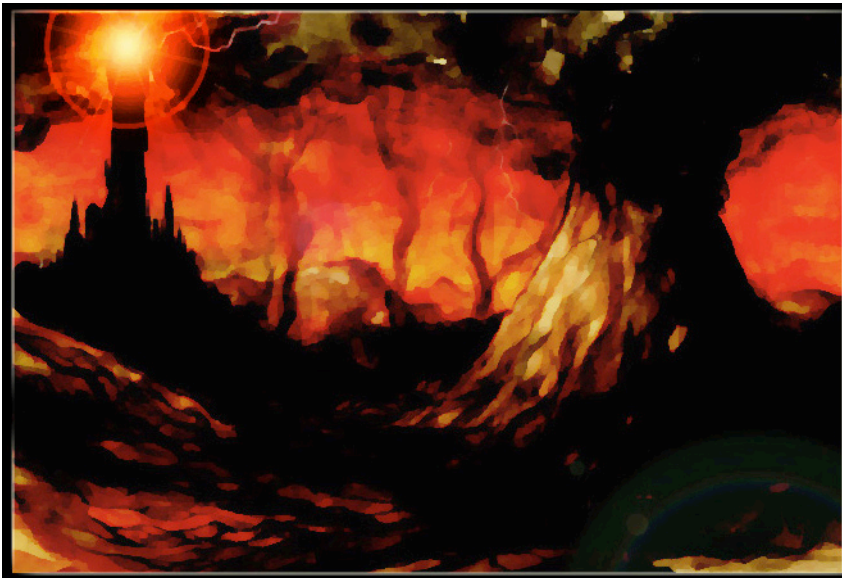
Phyla: Encapsulation

- Comparing Classes by Name
- Data Leaking Between Users
- Leftover Debug Code
- Mobile Code: Object Hijack
- Mobile Code: Use of Inner Class
- Mobile Code: Non-Final Public Field
- Private Array-Typed Field Returned from a Public Method
- Public Data Assigned to Private Array-Typed Field
- System Information Leak
- Trust Boundary Violation



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*. Environment





Bonus pernicious kingdom

■ Environment

- Everything that is outside of source code but is still critical to security
- ASP .NET Misconfiguration: Password in Configuration File, Insecure Compiler Optimization, and others
- Example: **ASP .NET Misconfiguration: Creating Debug Binary** phylum

```
<configuration>                                     // Debug binary
    <compilation debug="true">
        ...
    </compilation>
    ...
</configuration>
```



CLASSPATH

- Modifying the CLASSPATH environment variable is the equivalent of modifying a Windows/Unix PATH
 - An attacker can construct classes with “value added” features that perform malicious acts
 - Classic example is theft of username/password
 - Involves duping a user into running attacker's code



Weak access permissions (J2EE)

- EJB method permissions should never grant access to the ANYONE role
- Indicates that access control for an application has not been carefully thought through
- Method permissions should always be restricted to the minimum set of roles that should be granted access



Weak access permissions - example

The following example illustrates the improper use of method access controls...

```
<ejb-jar>
  ...
  <assembly-descriptor>
    <method-permission>
      <role-name>ANYONE</role-name>
      <method>
        <ejb-name>SomeBean</ejb-name>
        <method-name>someMethod</method-name>
      </method-permission>
    </assembly-descriptor>
  ...
</ejb-jar>
```



Phyla: Environment

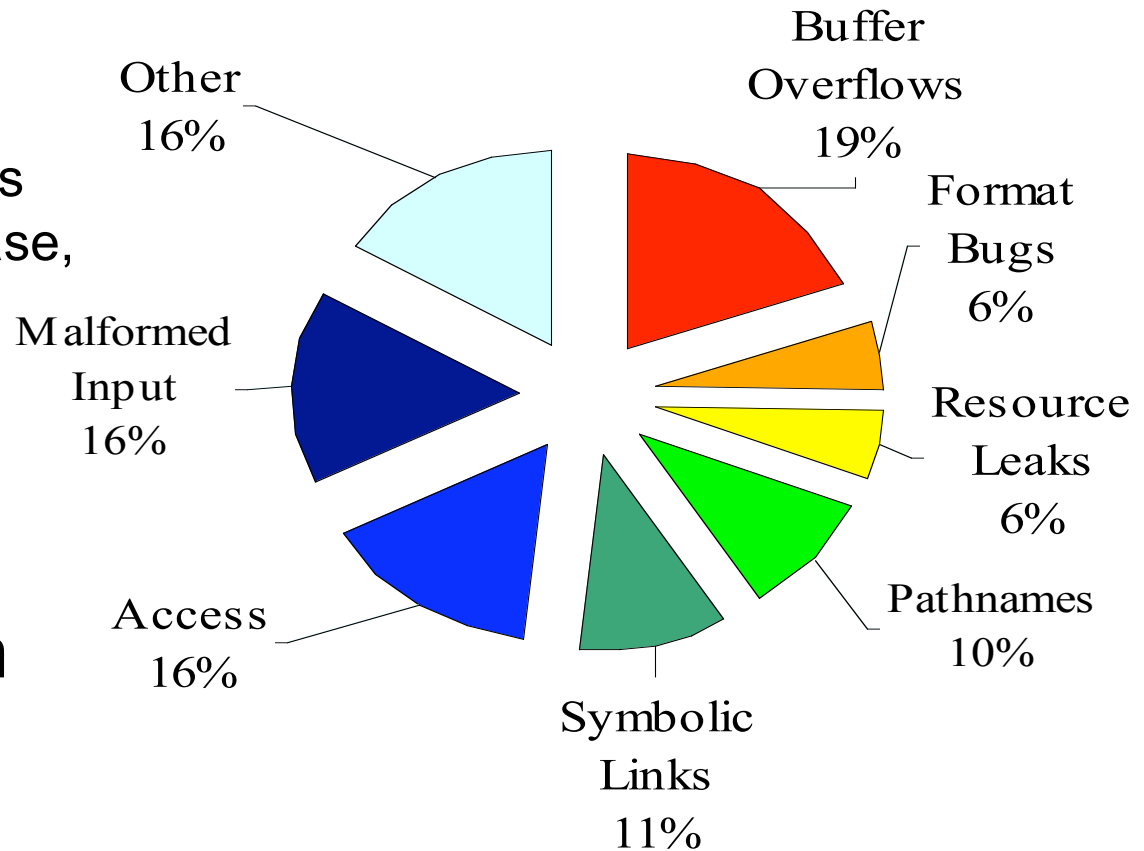
- **ASP .NET Misconfiguration: Creating Debug Binary**
- **ASP .NET Misconfiguration: Missing Custom Error Handling**
- **ASP .NET Misconfiguration: Password in Configuration File**
- **Insecure Compiler Optimization**
- **J2EE Misconfiguration: Insecure Transport**
- **J2EE Misconfiguration: Insufficient Session-ID Length**
- **J2EE Misconfiguration: Missing Error Handling**
- **J2EE Misconfiguration: Unsafe Bean Declaration**
- **J2EE Misconfiguration: Weak Access Permissions**

Static Analysis and Code Review



Reported flaws in
Common Vulnerabilities
and Exposures Database,
Jan-Sep 2001.

**56 % of CVE
vulnerabilities
could have been
detected with
straightforward
static analyses!**



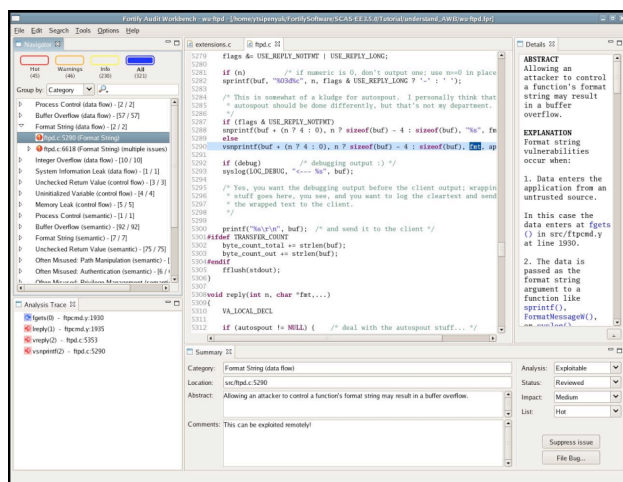
[Evans & Larochelle, IEEE Software, Jan 2002]



Touchpoint: code review (with a tool)

- Code review is a necessary evil
- Better coding practices make the job easier
- Automated tools help catch silly errors
 - Fortify/SCA (Cigital rules)
- Implementation errors do matter
 - Buffer overflows can be uncovered with static analysis
 - Static analysis
 - C/C++
 - Java
 - .NET
 - PSQL

- Tracing back from vulnerable location to input is critical





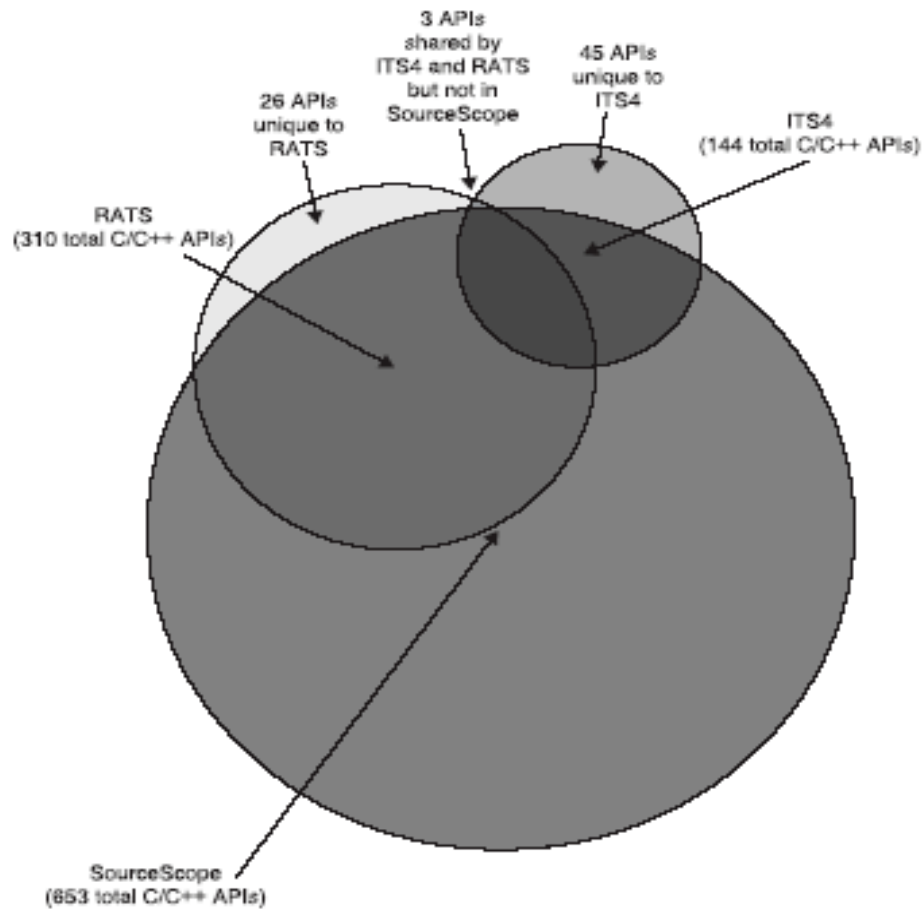
Code scanning tools

- Early static analysis tools (tokenizers)
 - ITS4
 - RATS
 - Flawfinder
- Modern tools (parsers)
 - Prefix
 - Fortify source code analysis suite
 - Ounce labs
 - Coverity
- The key is encapsulated know-how





Bug space coverage and early tools





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Fortify Source Code Analysis

- Integrated data flow analysis
- Broad platform support
- A comprehensive set of secure coding rules
 - Capability to add your own rules
- Proven large scale deployability

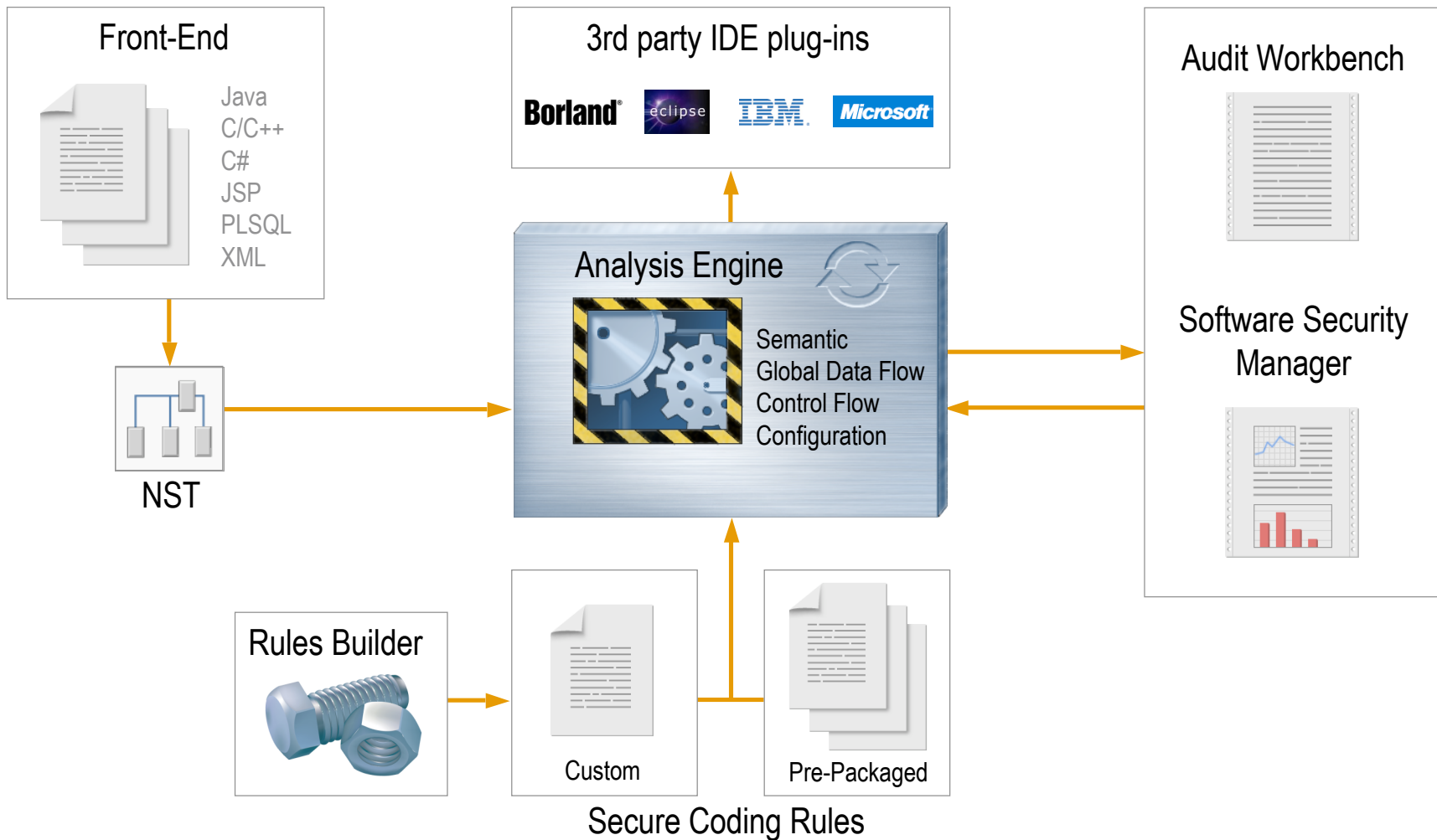
FORTIFY
SOFTWARE



Commercially viable, accurate and effective analysis



Fortify architecture





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Comprehensive secure coding rules

- Secure coding rulepacks based on the seven kingdoms
- Continuously updating and improving rulepacks
- Fortify Rules Builder allows you to further extend rulepacks to meet individualized needs
- Advanced context sensitive guidance inside in the IDE
- Intellectual property based on ten years of Cigital work
- see vulncat.fortifysoftware.com

The single largest compilation of secure coding techniques and guidance ever written



Next steps



Software security critical lessons

- Software security is more than a set of security functions
 - Not magic crypto fairy dust
 - Not silver-bullet security mechanisms
 - Not application of very simple tools
- Non-functional aspects of design are essential
- Security is an emergent property of the entire system (just like quality)
- Breaking stuff is important
- **To end up with secure software, deep integration with the SDLC is necessary**



Bottom up software security actions

- A few relatively simple things can make a tangible difference and can help you get started with software security

- Build checklists and use them
 - Sun's SAG checklist
<http://www.securecoding.org/companion/checklists/SAG/>
- Begin to develop a resource set (e.g., portal)
- Start small with simple architectural risk analyses (think Smurfware)
- Don't forget to include business-case justifications
- Use code scanning tools



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Top down software security actions

- Think of the problem as an evolutionary approach
- Chart out a strategic course of action to get where you want to be
 - Have a gap analysis performed
 - Make achievable, realistic milestones
 - Think about metrics for success
- Use outside help if you need it (Cigital)



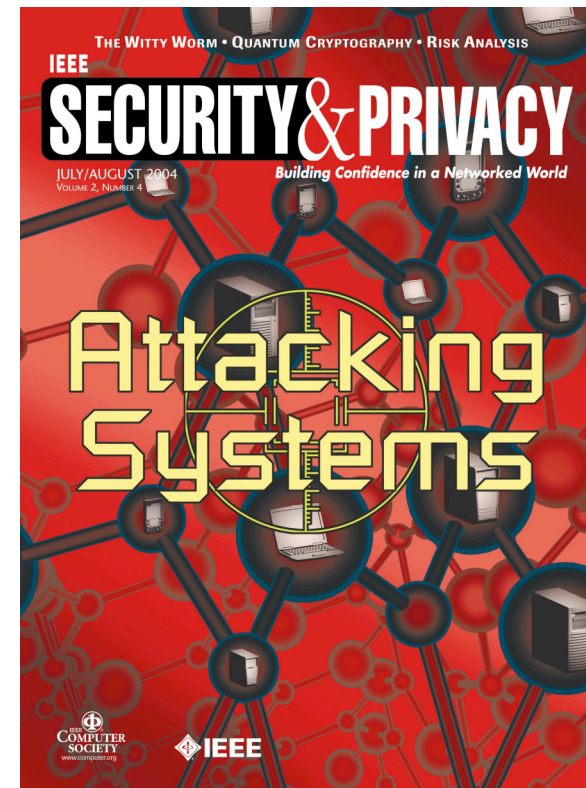


IEEE Security & Privacy Magazine

- Monthly Department on Software Security Best Practices called “Building Security In”

Table 1. Building Security In articles.

TITLE	KEY	AUTHOR	IEEE SECURITY & PRIVACY CITATION	URL: WWW.CIGITAL.COM PAPERS/DOWNLOAD/
Software Security	BSI1	Gary McGraw	2(2):80–83	bsi1-swsec.pdf
Misuse and Abuse Cases: Getting Past the Positive	BSI2	Paco Hope, Gary McGraw, Annie Anton	2(3):32–34	bsi2-misuse.pdf
Risk Analysis in Software Design	BSI3	Denis Verdon, Gary McGraw	2(4):79–84	bsi3-risk.pdf
Software Security Testing	BSI4	Bruce Potter, Gary McGraw	2(5):81–85	bsi4-testing.pdf
Static Analysis for Security	BSI5	Brian Chess, Gary McGraw	2(6):76–79	bsi5-static.pdf
Software Penetration Testing	BSI6	Brad Arkin, Scott Stender, Gary McGraw	3(1):84–87	bsi6-pentest.pdf
Knowledge for Software Security	BSI7	Sean Barnum, Gary McGraw	3(2):74–78	bsi7-knowledge.pdf
Adopting a Software Security Improvement Program	BSI8	Dan Taylor, Gary McGraw	3(3):88–91	bsi8-program.pdf
A Portal for Software Security	BSI9	Nancy R. Mead and Gary McGraw	3(4):75–79	bsi9-portal.pdf
Bridging the Gap between Software Development and Information Security	BSI10	Kenneth R. van Wyk and Gary McGraw	3(5):75–79	bsi10-bridge.pdf





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- See the Addison-Wesley Software Security series
- Send e-mail: gem@cigital.com
Ken@KRvW.com
- <http://www.cigital.com>
- <http://www.krvw.com>

“So now, when we face a choice between adding features and resolving security issues, we need to choose security.”

-Bill Gates



For more

